



Fundamentals of Digital Systems Design

Course website:

<http://www.liacs.leidenuniv.nl/~stefanovtp/courses/DITE/index.html>

Main Lecturer: Todor Stefanov (t.p.stefanov@liacs.leidenuniv.nl)

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Course Organization: Structure and Rules

- **On-campus Lectures**
 - your attendance is **strongly** recommended
- **On-campus Hands-on tutorials with practical assignments**
 - your attendance is **strongly** encouraged
- **On-campus and At-home Design Project**
 - you **must** make and turn-in the design project
- **Homeworks**
 - doing homeworks is **strongly** recommended
 - excellent material to practice for the exams!

Important:

- All info needed for the course will be on the course website
- Check the website regularly for instructions
- Check regularly Brightspace and your email for announcements



Course Organization: On-campus Lectures

■ Goal

- To learn the fundamentals of digital systems design
- Theoretical Foundations
- Basic Digital Circuits and Building Blocks
- Design Algorithms and Procedures
- The basics of Memory Design
- The basics of Micro-Processor Design

■ Format

- **On-campus presentations** supported by PowerPoint slides
- Slides available after each lecture on the course website

■ Literature

- M. Morris Mano and Charles R. Kime,
"Logic and Computer Design Fundamentals", 3rd or 4th edition

Important:

- Dedicated 2 hours for questions after each lecture on Mondays!

Lectures Schedule and Location

Date	Time	Location
04 Sept (Monday)	15:15 - 17:00	Lecture Hall - C1
06 Sept (Wednesday)	15:15 - 17:00	Lecture Hall - C1
11 Sept (Monday)	15:15 - 17:00	Lecture Hall - C1
13 Sept (Wednesday)	15:15 - 17:00	Lecture Hall - C1
25 Sept (Monday)	15:15 - 17:00	Lecture Hall - C1
27 Sept (Wednesday)	15:15 - 17:00	Lecture Hall - C1
09 Oct (Monday)	15:15 - 17:00	Lecture Hall - C1
16 Oct (Monday)	15:15 - 17:00	Lecture Hall - C1
30 Oct (Monday)	15:15 - 17:00	Lecture Hall - C1
06 Nov (Monday)	15:15 - 17:00	Lecture Hall - C1
13 Nov (Monday)	15:15 - 17:00	Lecture Hall - C1
20 Nov (Monday)	15:15 - 17:00	Lecture Hall - C1
27 Nov (Monday)	15:15 - 17:00	Lecture Hall - C1
04 Dec (Monday)	15:15 - 17:00	Lecture Hall - C1
11 Dec (Monday)	15:15 - 17:00	Lecture Hall - C1



Course Organization: On-campus Tutorials

■ Goal

- To learn how to solve simple digital design problems in practice
- To learn how to use software tools to design digital circuits

■ Format

- Pen-and-paper practical assignments (simple digital circuit designs)
- Design and simulation of digital circuits using software tools

■ Tools

- Xilinx ISE for digital circuit design and simulation
- Go to the website of the course to see options to access Xilinx ISE!

■ Timetable and Location

Period	Day	Time	Place
11 Oct – 08 Nov	Every Wednesday	15:15 - 17:00	Rooms 302-304, 303, 306-308, 307, 309 in Snellius



Course Organization: Design Project

■ Goal

- To apply the knowledge gained at lectures and hands-on tutorials
- To design a simple 4-bit Micro-Processor

■ Format

- Project specification will be given in beginning of November
- Project tutorials with the tutors on Wednesdays
 - to work on the project on-campus
 - to get advice on how to design the micro-processor
 - to ask questions
- You have to spend extra time at home or in our computer labs if
 - project tutorials time is not sufficient to accomplish the project

■ Timetable and Location

Period	Day	Time	Place
15 Nov - 13 Dec	Every Wednesday	15:15 - 17:00	Rooms 302-304, 303, 306-308, 307, 309 in Snellius



Course Organization: Homeworks

■ Goals

- To understand better the material given at lectures
- To study alone additional material not given at lectures
- To prepare and practice for exams!

■ Format

- Several homeworks will be given
- Homework instructions posted on website every Wednesday
- Try to do homework alone
 - within 7 days after the instructions are posted
- Answers to homeworks will be posted on the course website
 - 7 days after the instructions

Important:

- Dedicated 2 hours for questions after each lecture on Mondays!

Course Organization: Exams and Grading

■ Exams

- One *Midterm exam* and one *Final exam* at the end of the course
- You **must** be present at these exams
- In case you fail at the final exam, there will be one Re-take exam

■ Timetable and Location

Exams	Date	Hour	Place
Midterm Exam	October 25, 2023	13:15 – 16:15	Lecture Hall - C4/5
Final Exam	January 26, 2024	13:15 – 16:15	Lecture Hall - C4/5

■ Grading

- You will receive 3 Grades: one for Final exam (G^{Fexam}), one for Project (G^{Project}), and one for Midterm exam (G^{Mexam}),
- All grades are important and will form your **Final Grade** as follows:

If ($G^{\text{Fexam}} \geq 6.0$ and $G^{\text{Project}} \geq 6.0$) Then

$$\text{Final Grade} = 0.4 * G^{\text{Fexam}} + 0.4 * G^{\text{Project}} + 0.2 * G^{\text{Mexam}}$$

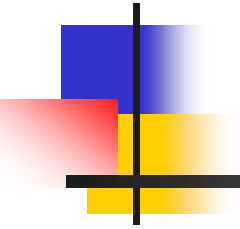
Else

$$\text{Final Grade} \leq 5$$



Digital Systems and Information

Part I





Overview

- Introduction to Digital Systems
 - Digital Systems (examples and general remarks)
 - Basic Digital System Structure (basic components and their function)
- Number Systems
 - Positional Number Systems (decimal, binary, octal, and hexadecimal)
 - Number Conversions (r-to-decimal, decimal-to-r, other conversions)
- Representations of Numbers in Digital Systems
 - Integer Numbers (unsigned and signed representations)
 - Real Numbers (floating-point representation)

Digital Systems: Example

- Modern Digital Personal Computers (PC)
 - The best known example of a digital system
 - Most striking properties are: **generality** and **flexibility**



Digital Systems: More Examples





Digital Systems: General Remarks

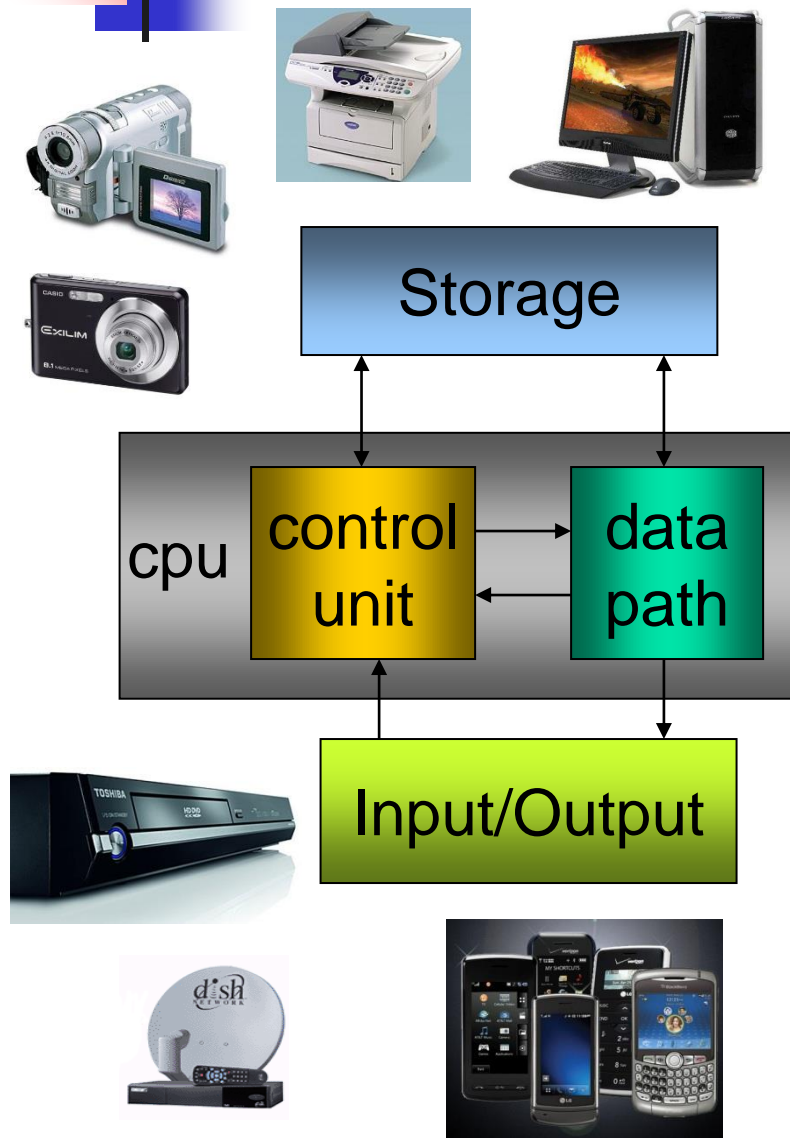
- Digital Systems **manipulate discrete elements/quantities** of information
- Discrete quantities of information emerge from:
 - the nature of the data being processed
 - the data may be discretized from continuous values
- Early computer systems used mainly for numeric computations
 - the discrete elements used were the digits
 - hence the term **digital computer/system** emerged
- In general, any system uses an **alphabet** (set of **symbols**) to represent information
 - The English language system uses an alphabet of 26 symbols (letters)
 - The decimal number system uses an alphabet of 10 symbols (digits)
- What about the alphabet of the Digital Systems?



The Digital Systems' Alphabet is Binary

- Digital Systems use only one alphabet with two symbols (digits) '0' and '1' (hence **binary alphabet**)
 - Binary digit is called a **bit**
 - Information is represented by groups of bits
- Why is a binary alphabet used?
 - Digital systems have a basic building block called **switch**
 - The switch can only be "**on**" or "**off**"
 - That is, two discrete values '0' and '1' can be distinguished
- Electric device, **transistor**, physically implements the switch
 - '0' and '1' are physically represented by voltage values called **LOW** and **HIGH**
 - "**on**" (**closed**) switch corresponds to '0' and is represented by **LOW** voltage value (between 0.0 and 1.0 Volt)
 - "**off**" (**open**) switch corresponds to '1' and is represented by **HIGH** voltage value (between 3.0 and 5.0 Volts)
 - More information will be given later in another lecture

Basic Digital System Structure



- **CPU**: Central Processing Unit
- **Data Path**: arithmetic and logic operations
- **Control Unit**: make sure that the sequence of data path operations is correct
- **Storage**: no memory = no system
- **Input/Output**: allow the system to interact with the outside world



Information Representation

- All information in Digital Systems is represented in binary form
- Information that is not binary is converted to binary before processed by Digital Systems
- Decimal numbers are expressed
 - in the binary number system or
 - by means of a binary code
- That is not too difficult because
 - All number systems have a similar formal representation
 - Thus, one number system can be converted into another
- Let us look into number systems and conversions



Number Systems

Number Systems are used in arithmetic to represent numbers by strings of digits. There are two types of systems:

- **Positional** number systems

- The meaning of each digit depends on its position in the number
- Example: Decimal number system
 - **585.5** is a decimal number in positional code
 - “**5** hundreds plus **8** tens plus **5** units plus **5** tenths”
 - The hundreds, tens, units, and tenths are powers of 10 implied by the position of the digits

$$\mathbf{585.5} = \mathbf{5} \times 10^2 + \mathbf{8} \times 10^1 + \mathbf{5} \times 10^0 + \mathbf{5} \times 10^{-1}$$

- Decimal number system is said to be of base or radix 10 because
 - it uses 10 distinct digits (0 – 9)
 - the digits are multiplied by powers of 10
- **Non-positional** number systems

- Old Roman numbers: for example, **XIX** equals to 19



Positional Number Systems

We can represent numbers in any number system with base r

■ Number in positional code

- $(A_{n-1}A_{n-2}\dots A_1A_0.A_{-1}A_{-2}\dots A_{-m+1}A_{-m})_r$
- r is the base (radix) of the system, $r \in \{2, 3, \dots, \infty\}$
- every digit $A_i \in \{0, 1, 2, \dots, r-1\}$, where $\{0, 1, 2, \dots, r-1\}$ is the digit set
- “.” is called the radix point
- A_{n-1} is referred to as the most significant digit
- A_{-m} is referred to as the least significant digit

■ Number in base r expressed as power series of r

- $A_{n-1}r^{n-1} + A_{n-2}r^{n-2} + \dots + A_1r^1 + A_0r^0 + A_{-1}r^{-1} + A_{-2}r^{-2} + \dots + A_{-m+1}r^{-m+1} + A_{-m}r^{-m}$

■ Example: a number in number system with base 5

- $(132.4)_5 = 1 \times 5^2 + 3 \times 5^1 + 2 \times 5^0 + 4 \times 5^{-1} = 25 + 15 + 2 + 0.8 = (42.8)_{10}$



Binary Number System

This is the system used for arithmetic in all digital computers

- Number in positional code

- $(b_{n-1}b_{n-2}\dots b_1b_0.b_{-1}b_{-2}\dots b_{-m+1}b_{-m})_r$
- $r = 2$ is the base of the binary system
- every digit $b_i \in \{0, 1\}$
- the digits b_i in a binary number are called **bits**
- b_{n-1} is referred to as the most significant bit (MSB)
- b_{-m} is referred to as the least significant bit (LSB)

- Number in base 2 expressed as power series of 2

- $b_{n-1}2^{n-1} + b_{n-2}2^{n-2} + \dots + b_12^1 + b_02^0 + b_{-1}2^{-1} + b_{-2}2^{-2} + \dots + b_{-m+1}2^{-m+1} + b_{-m}2^{-m}$

- Example: a number in the binary number system

- $(1011.01)_2 = 1 \times 2^3 + 0 \times 2^2 + 1 \times 2^1 + 1 \times 2^0 + 0 \times 2^{-1} + 1 \times 2^{-2} = 8 + 2 + 1 + 0.25 = (11.25)_{10}$

Burn this table into your memory

Power of Two

n	2^n	n	2^n	n	2^n
0	1	5	32	10	1024
1	2	6	64	11	2048
2	4	7	128	12	4096
3	8	8	256	13	8192
4	16	9	512	14	16384

$2^{10} = 1024$ **is 1K**
(kilo)

2^{20} **is 1M**
(mega)

2^{30} **is 1G**
(giga)

Other Useful Number Systems

Octal (base-8) and Hexadecimal (base-16) number systems

- useful for representing binary quantities indirectly because
- their bases are powers of two
- have more compact representations of binary quantities

■ Octal number system

- $(o_{n-1}o_{n-2}\dots o_1o_0.o_{-1}o_{-2}\dots o_{-m+1}o_{-m})_8$
- every digit $o_i \in \{0, 1, 2, 3, 4, 5, 6, 7\}$.
- $o_{n-1}8^{n-1} + o_{n-2}8^{n-2} + \dots + o_18^1 + o_08^0 + o_{-1}8^{-1} + o_{-2}8^{-2} + \dots + o_{-m+1}8^{-m+1} + o_{-m}8^{-m}$
- $(127.4)_8 = 1 \times 8^2 + 2 \times 8^1 + 7 \times 8^0 + 4 \times 8^{-1} = (87.5)_{10} = (001\ 010\ 111.100)_2$

■ Hexadecimal number system

- $(h_{n-1}h_{n-2}\dots h_1h_0.h_{-1}h_{-2}\dots h_{-m+1}h_{-m})_{16}$
- every digit $h_i \in \{0, 1, 2, 3, 4, 5, 6, 7, 8, 9, A, B, C, D, E, F\}$.
- $h_{n-1}16^{n-1} + h_{n-2}16^{n-2} + \dots + h_116^1 + h_016^0 + h_{-1}16^{-1} + h_{-2}16^{-2} + \dots + h_{-m+1}16^{-m+1} + h_{-m}16^{-m}$
- $(B6F.4)_{16} = 11 \times 16^2 + 6 \times 16^1 + 15 \times 16^0 + 4 \times 16^{-1} = (2927.25)_{10} = (1011\ 0110\ 1111.0100)_2$



Another Important Table

Decimal (base 10)	Binary (base 2)	Octal (base 8)	Hex (base 16)
0	0000	0	0
1	0001	1	1
2	0010	2	2
3	0011	3	3
4	0100	4	4
5	0101	5	5
6	0110	6	6
7	0111	7	7
8	1000	10	8
9	1001	11	9
10	1010	12	A
11	1011	13	B
12	1100	14	C
13	1101	15	D
14	1110	16	E
15	1111	17	F

Conversion from base r to Decimal

To convert a number in base r to decimal number (base 10) do:

- expand the number in power series of r and
- add all the terms as shown below

$$(A_{n-1}A_{n-2}\dots A_1A_0.A_{-1}A_{-2}\dots A_{-m+1}A_{-m})_r =$$

$$A_{n-1}r^{n-1} + A_{n-2}r^{n-2} + \dots + A_1r^1 + A_0r^0 + A_{-1}r^{-1} + A_{-2}r^{-2} + \dots + A_{-m+1}r^{-m+1} + A_{-m}r^{-m}$$

- Example of converting Binary (base 2) to Decimal (base 10):

$$(1011.01)_2 = 1x2^3 + 0x2^2 + 1x2^1 + 1x2^0 + 0x2^{-1} + 1x2^{-2} = 8 + 2 + 1 + 0.25 = (11.25)_{10}$$

- Example of converting number in base 5 to Decimal (base 10):

$$(132.4)_5 = 1x5^2 + 3x5^1 + 2x5^0 + 4x5^{-1} = 25 + 15 + 2 + 0.8 = (42.8)_{10}$$

- Example of converting Octal (base 8) to Decimal (base 10):

$$(127.4)_8 = 1x8^2 + 2x8^1 + 7x8^0 + 4x8^{-1} = (87.5)_{10}$$

- Example of converting Hexadecimal (base 16) to Decimal (base 10):

$$(B6F.4)_{16} = 11x16^2 + 6x16^1 + 15x16^0 + 4x16^{-1} = (2927.25)_{10}$$

Conversion from Decimal to base r

The conversion is done as follows:

- 1) Separate the number into **integer part** and **fraction part**
 - the two parts must be converted differently
- 2) Converting the integer part
 - divide the integer part and all successive quotients by r and
 - accumulate the remainders
- 3) Converting the fraction part
 - multiply the fractional parts by r and accumulate integers

■ Example of converting Decimal (base 10) to Binary (base 2): **$(41.6875)_{10}$**

Converting the integer part **$(41)_{10}$** :

$41/2 = 20 + 1/2$	Remainder = 1	↑	LSB
$20/2 = 10 + 0/2$	0		
$10/2 = 5 + 0/2$	0		
$5/2 = 2 + 1/2$	1		
$2/2 = 1 + 0/2$	0		
$1/2 = 0 + 1/2$	1	↓	MSB

$$(41)_{10} = (101001)_2$$

Converting the fraction part **$(0.6875)_{10}$** :

$0.6875 \times 2 = 1.3750$	Integer = 1	↓	MSB
$0.3750 \times 2 = 0.7500$	0		
$0.7500 \times 2 = 1.5000$	1		
$0.5000 \times 2 = 1.0000$	1	↓	LSB

$$(0.6875)_{10} = (0.1011)_2$$

$$(41.6875)_{10} = (101001.1011)_2$$



Other Conversions

- Binary to Octal or Hexadecimal: grouping bits starting from the radix point
 - $(1101010.01)_2$ to Octal (groups of 3): $(001|101|010.010|)_2 = (152.2)_8$
 - $(1101010.01)_2$ to Hex (groups of 4): $(0110|1010.0100|)_2 = (6A.4)_{16}$
- Octal to Binary: convert each digit to binary using 3 bits
 - $(475.2)_8 = (100\ 111\ 101.\ 010)_2$
- Hexadecimal to Binary: convert each digit to binary using 4 bits
 - $(7A5F.C)_{16} = (0111\ 1010\ 0101\ 1111.\ 1100)_2 = (111101001011111.11)_2$
- Hexadecimal to Octal
 - Hexadecimal \rightarrow Binary \rightarrow Octal
- Octal to Hexadecimal
 - Octal \rightarrow Binary \rightarrow Hexadecimal

Range of Numbers in base r

- Assume numbers represented by n digits for the integer part and m digits for the fraction part -- $(A_{n-1}A_{n-2}\dots A_1A_0.A_{-1}A_{-2}\dots A_{-m+1}A_{-m})_r$
- What is the smallest and the largest number that can be represented?
- The smallest integer number is **0** and the largest is $(r-1)r^{n-1} + (r-1)r^{n-2} + \dots + (r-1)r^1 + (r-1)r^0 = r^n - 1$, i.e., **the range is from 0 to $r^n - 1$**
- The smallest fraction number is **0.0** and the largest is $(r-1)r^{-1} + (r-1)r^{-2} + \dots + (r-1)r^{-m+1} + (r-1)r^{-m} = 1 - r^{-m}$, i.e., **the range is from 0.0 to $1 - r^{-m}$**
- The range of numbers is **from 0.0 to $r^n - r^{-m}$**
 - The range of numbers in base r depends on the number of digits used to represent the numbers
- Examples:
 - Largest 3-digit integer decimal (base 10) number is $10^3 - 1 = 1000 - 1 = 999$
 - Largest 8-digit integer binary (base 2) number is $(11111111)_2$, i.e., $2^8 - 1 = 255$
 - Largest 5-digit decimal (base 10) fraction is $1 - 10^{-5} = 1 - 0.00001 = 0.99999$
 - Largest 16-digit binary (base 2) fraction is $1 - 2^{-16} = 0.9999847412$
- What about the range of negative numbers?

Representations of Numbers in Digital Systems (1)

- Numbers are represented in binary format as strings of bits
 - **Bit** is the smallest binary quantity with a value of 0 or 1
 - **Byte** is a string (sequence) of 8 bits
 - **Word** is a string (sequence) of n bits ($n > 8$)
 - Examples → **bit**: 1 **byte**: 01101111 **16-bit word**: 11110100 10001010
- Positive Integer Numbers
 - Can be represented as **Unsigned Binary Numbers** using a string of n bits
 - **Magnitude representation** – number X in binary having n bits
 - Example: **00001001** (represents integer number 9 using 8 bits)
 - **2's Complement representation** – Given number X in binary having n bits, the 2's complement of X is defined as $2^n - X$.
 - Example: **11110111** (2's complement of integer number 9)
 - **1's Complement representation** - Given number X in binary having n bits, the 1's complement of X is defined as $(2^n - 1) - X$.
 - Example: **11110110** (1's complement of integer number 9).

Representations of Numbers in Digital Systems (2)

- Positive and Negative Integer Numbers
 - Can be represented as **Signed Binary Numbers** using a string of n bits
 - The most significant bit is interpreted as a sign bit
 - For positive numbers the sign bit is 0
 - For negative numbers the sign bit is 1
 - Signed-Magnitude representation
 - Example: **0|0001001** (represents integer number +9 using 8 bits)
 - Example: **1|0001001** (represents integer number -9 using 8 bits)
 - Signed-2's *Complement representation*
 - Example: **0|0001001** (signed-2's complement of integer number +9)
 - Example: **1|1110111** (signed-2's complement of integer number -9)
 - Signed-1's *Complement representation*
 - Example: **0|0001001** (signed-1's complement of integer number +9)
 - Example: **1|1110110** (signed-1's complement of integer number -9)
- How do we get the signed complements?

Representations of Numbers in Digital Systems (3)

- Real Numbers
 - Can be represented as **fixed-point** or **floating-point** numbers
- We will look closer to **floating-point** numbers
 - Floating-point representation is similar to scientific notation
 - It consists of two parts – mantissa **M** and exponent **E**
 - Floating-point number **F** represented by the pair **(M,E)** has the value

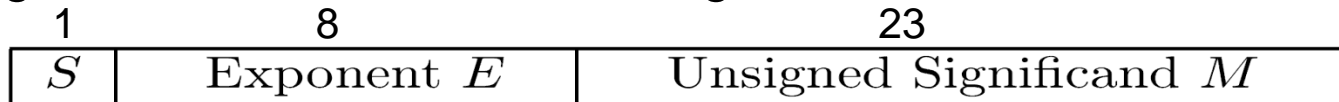
$$F = M \times \beta^E \quad (\beta - \text{base of exponent})$$

- Example: the decimal number **-179.75** as a floating-point number is **-0.17975 x 10⁺³**
 - Mantissa **M** = -0.17975 – represents the number as a fraction
 - Exponent **E** = +3 – designates the position of the radix point
 - Base **β** = 10 – the number is decimal
- Floating-point representation is not unique
 - -179.75 = -0.17975 x 10⁺³ = -0.0017975 x 10⁺⁵ .
 - -0.17975 x 10⁺³ is called **normalized** form because the most significant digit of the fraction is nonzero.

Representations of Numbers in Digital Systems (4)

- IEEE standard **single-precision floating-point representation**

- Using 32-bit word with the following format:



- The value of the number is: $(-1)^S \times (1.M)_2 \times 2^{E-127}$
- The largest and smallest positive numbers are:
 - $(-1)^0 \times (1.111111111111111111111111111111)_2 \times 2^{254-127} = +(1+1-2^{-23}) \times 2^{127}$
 - $(-1)^0 \times (1.000000000000000000000000000000)_2 \times 2^{1-127} = + 2^{-126}$
- Special cases:
 - $E = 255$ and $M = 0$, the number represents plus or minus infinity.
 - $E = 255$ and $M \neq 0$, not a valid number (signify invalid operations).
 - $E = 0$ and $M = 0$, the number denotes plus or minus zero.
 - $E = 0$ and $M \neq 0$, the number is said to be de-normalized.

- Example: $(-179.75)_{10} = (-10110011.11)_2 = (-1.011001111 \times 2^7)_2$

