



Universiteit Leiden

Applicability of Loop Recombination in Ciliates using the Breakpoint Graph

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and Michael Muskulus

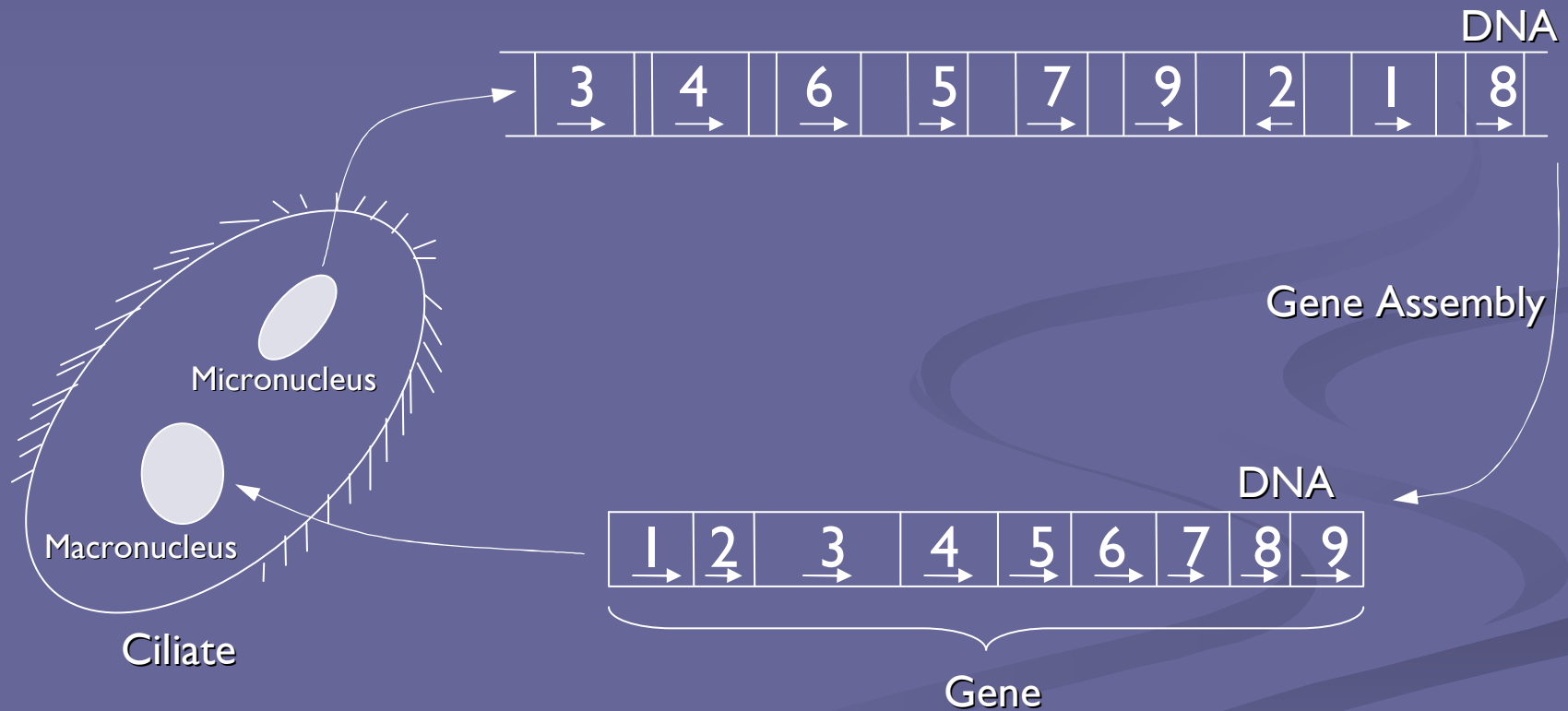
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Overview

- Brief overview of gene assembly in ciliates.
- Brief overview of a formal model.
- Recall reduction graph.
- Introduce graph \tilde{G} on top of G 's reduction graph, and show its uses.

Gene Assembly in Ciliates

Agene of the *Sterkiella nova*

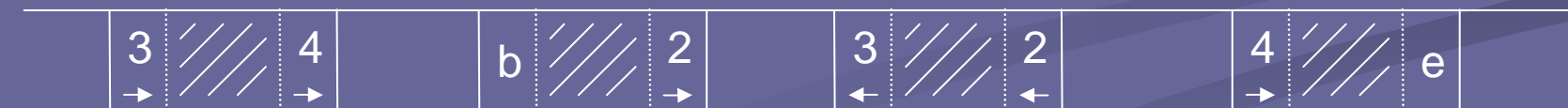


GeneAssembly

- The gene assembly process is done using molecular operations.
- With the aid of *pointers* these operations `know' how these parts need to be glued together.



Actually



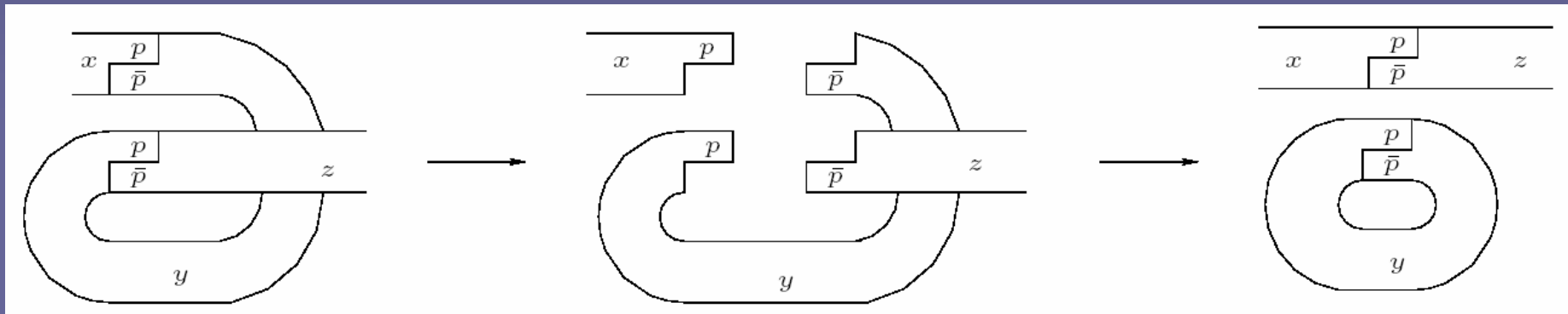
GeneAssembly

The gene assembly process is accomplished using three molecular operations:

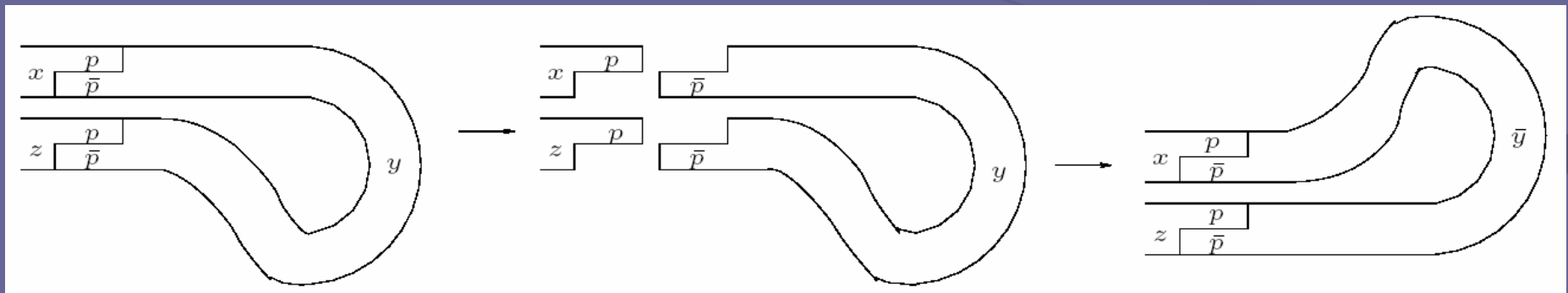
- Loop recombination
- Hairpin recombination
- Double-loop recombination

GeneAssembly

- Looprecombination

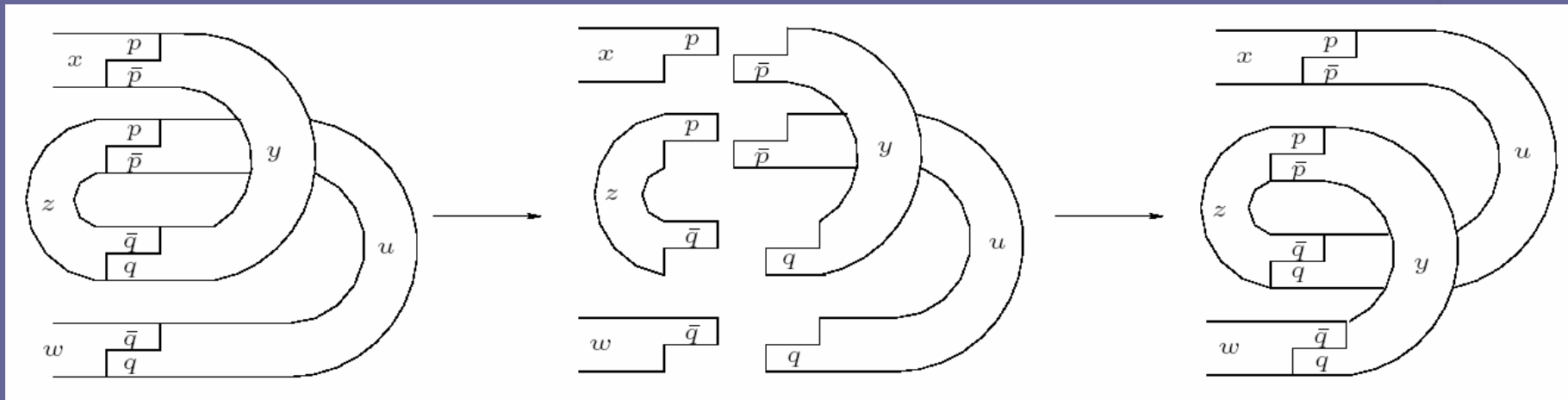


- Hairpinrecombination



GeneAssembly

- Double-loop recombination



GeneAssembly

The process is irreversible:

when a molecular operation is applied on a pointer, then this pointer cannot be used again.

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Modeling Gene Assembly



SPRS

342 $\overline{3}$ $\overline{2}$ 4 legalstring

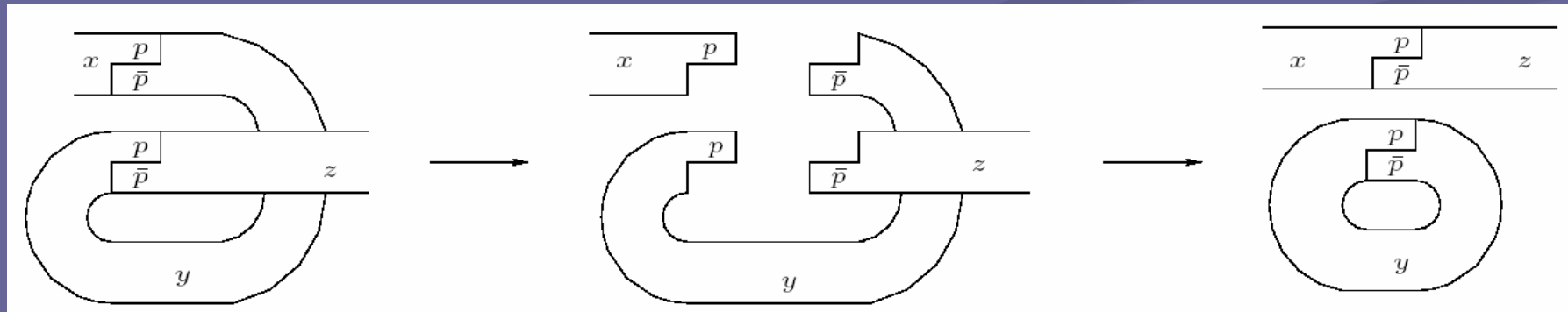
$\Pi = \{2, \overline{2}, 3, \overline{3}, \dots\}$.

Modeling Gene Assembly

Loop recombination - String negation rule

For $p \in \Pi$ and $x, z \in \Pi^*$:

$$\text{snr}_p(xppz) = xz$$



e.g., $\text{snr}_3(45\bar{3}\bar{3}45) = 4545$.

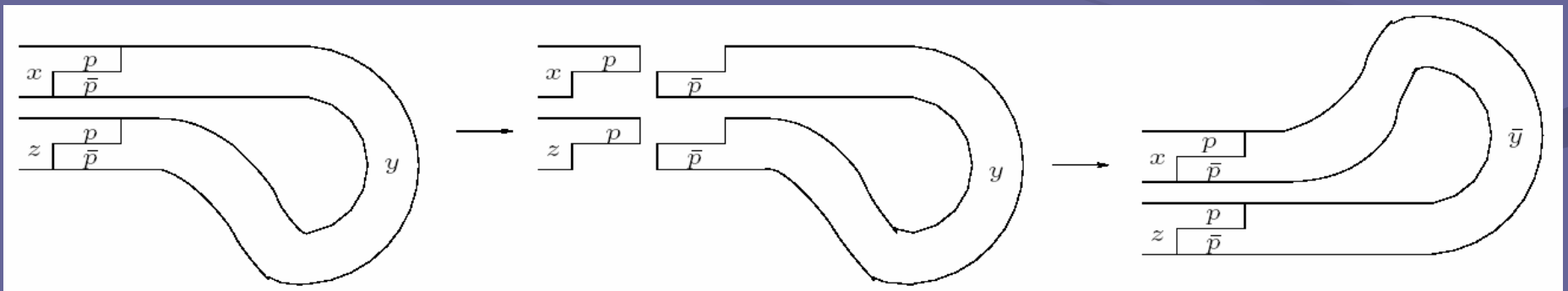
Modeling Gene Assembly

Hairpin recombination - String positive rule

$$\text{spr}_p(xpy\bar{p}z) = x\bar{y}z$$

where $\bar{\bar{p}} = p$

$$u = x_1x_2 \cdots x_n \Rightarrow \bar{u} = \bar{x}_n\bar{x}_{n-1} \cdots \bar{x}_1$$



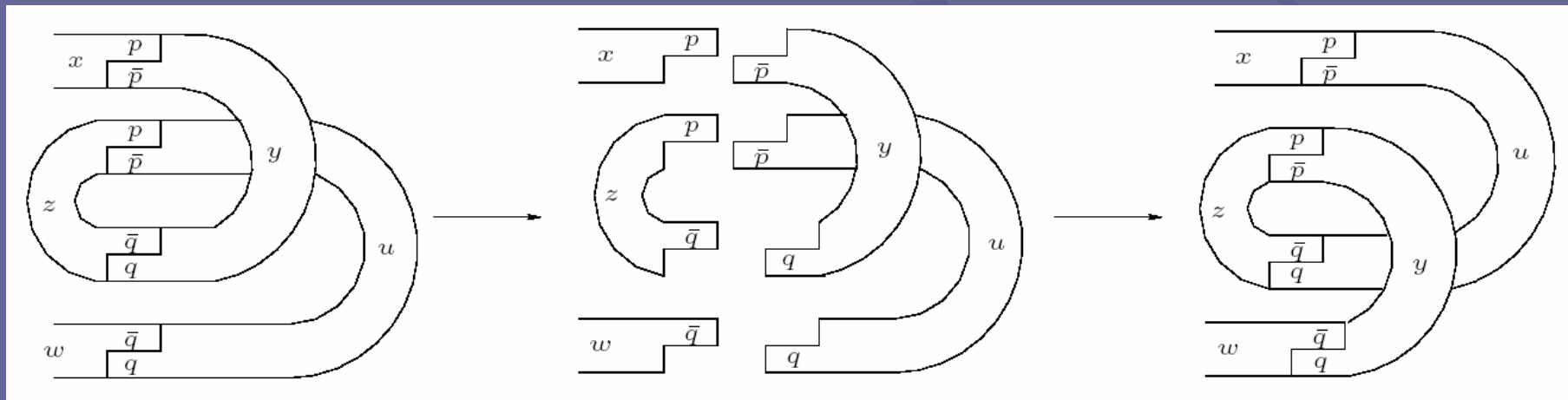
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Modeling Gene Assembly

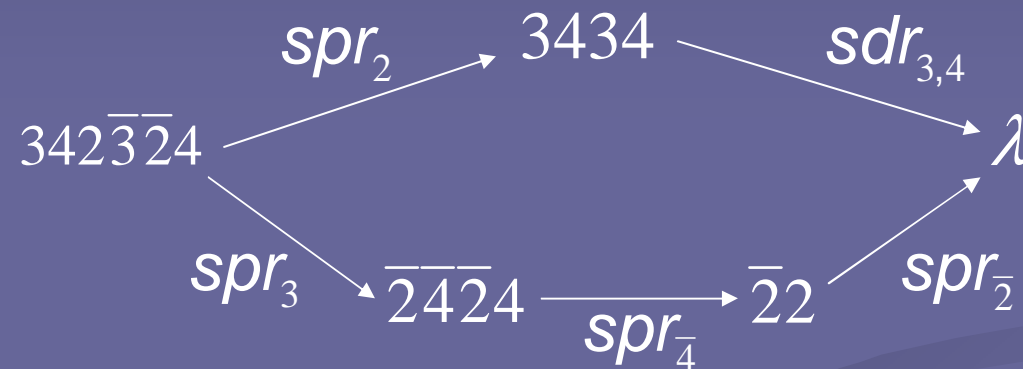
Double-loop recombination - String doubling rule

$$\text{sdr}_{p,q}(xpyqzpuqw) = xuzyw$$



Modeling Gene Assembly

- Successful reductions:



- λ correspondstoasuccessfulassembledgene.
- Non-deterministicprocess.
- Nodeadlocks.
- Finiteprocess.Pointersareremovedineverystep.



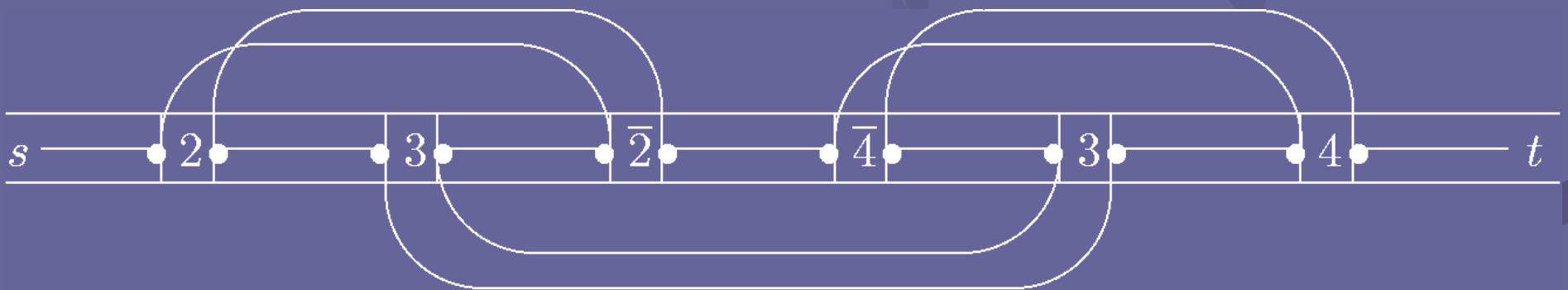
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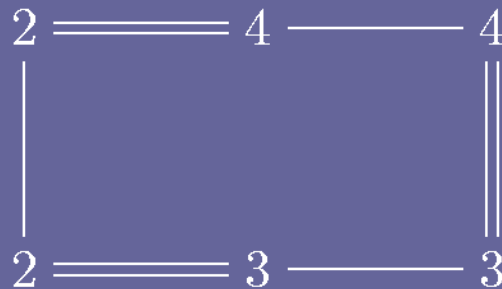
ReductionGraph



Concept of breakpoint graph:
reality-and-desire



Reduction Graph Example



- Double edges are the reality edges, single edges the desire edges
- There is a cyclic component.
- This is the less general version of reduction graph.

Known result

Let u be a legal string, and let N be the number of cyclic components in $R(u)$. Then every successful reduction of u has exactly N string negative rules.

Complexity: $O(|u|)$

Example

Let $u = 23\overline{24}34$ be a legal string. Then $R(u)$ has one cyclic component.
 Thus, every successful reduction of u has exactly one string negative rule.

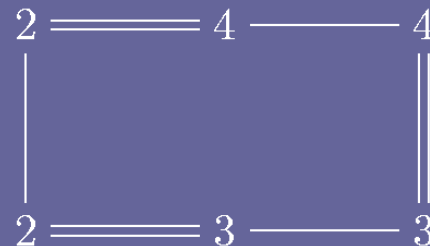
Successful reductions :

$snr_2 spr_3 spr_4, \quad s \text{ --- } 2 \text{ --- } 2 \text{ --- } 3 \text{ --- } 3 \text{ --- } 4 \text{ --- } 4 \text{ --- } t$

$snr_3 spr_2 spr_4,$

$snr_3 spr_4 spr_2,$

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Overview

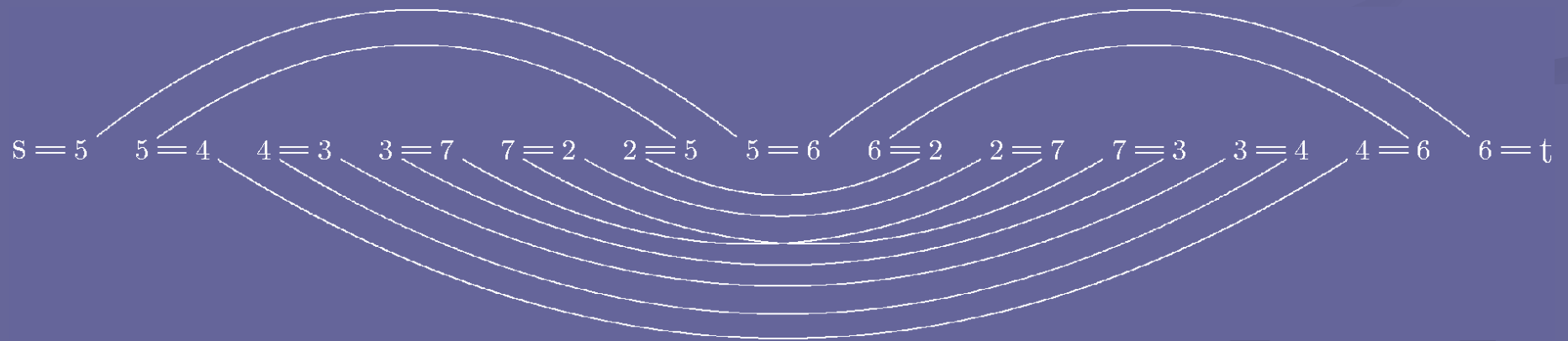
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- Brief overview of a formal model.
- Recall reduction graph.
- ***Introduce graph 'ontop of' reduction graph, and show its uses.***

Motivation

- Now we know *how many* Snr rules are needed, can we also determine *on which pointers* these Snr rules are applied in successful reductions?
- Previous example: Snr domains are {2}, {3}, {4}.
- If so, is this information retained in the reduction graph?

More complicated Example

$u = 54372562\bar{7}346$



More complicated Example

- Which is:

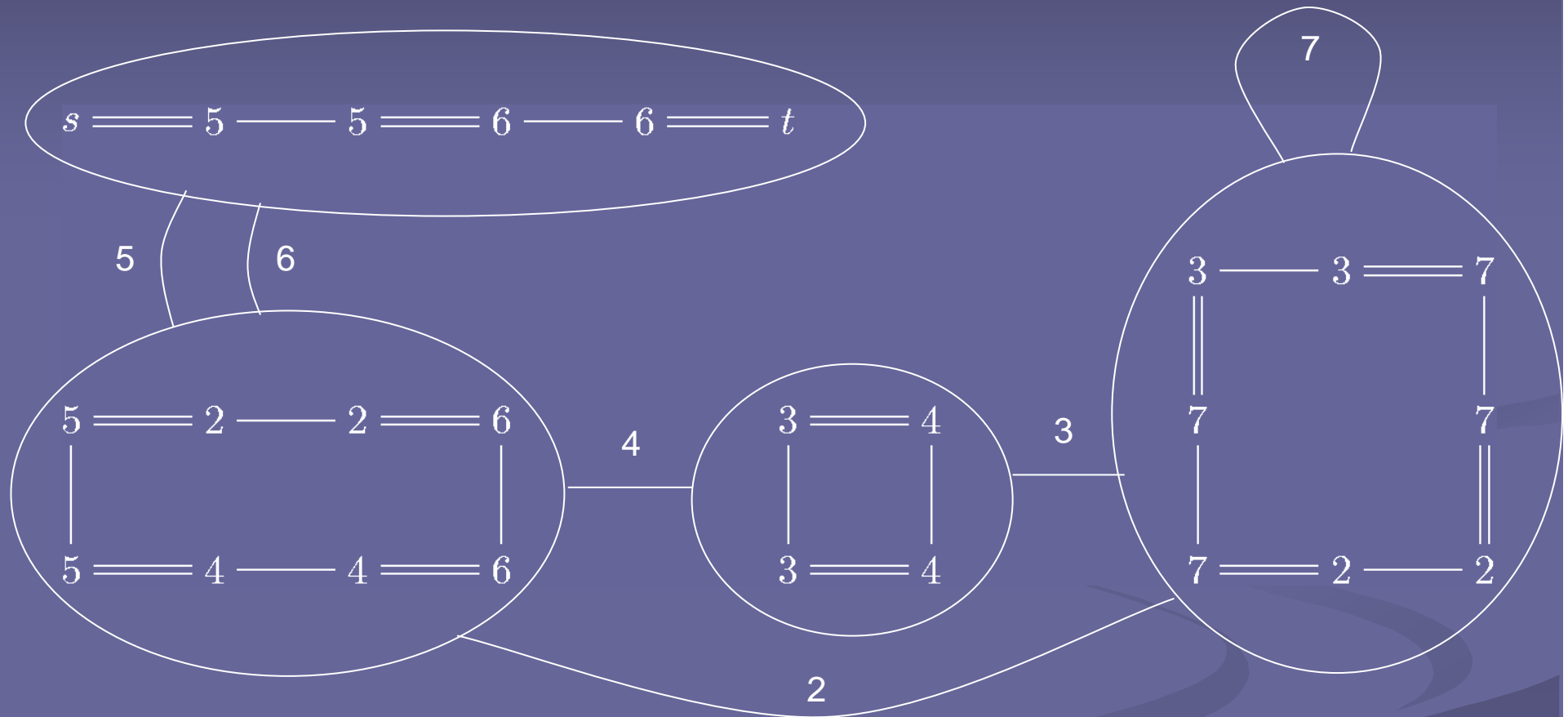
$$s \equiv 5 \text{ --- } 5 \equiv 6 \text{ --- } 6 \equiv t$$

$$\begin{array}{ccccc} 5 & \equiv & 2 & \text{---} & 2 \equiv 6 \\ | & & & & | \\ 5 & \equiv & 4 & \text{---} & 4 \equiv 6 \end{array}$$

$$\begin{array}{cc} 3 & \equiv & 4 \\ | & & | \\ 3 & \equiv & 4 \end{array}$$

$$\begin{array}{ccc} 3 & \text{---} & 3 \equiv 7 \\ || & & | \\ 7 & & 7 \\ | & & || \\ 7 \equiv 2 & \text{---} & 2 \end{array}$$

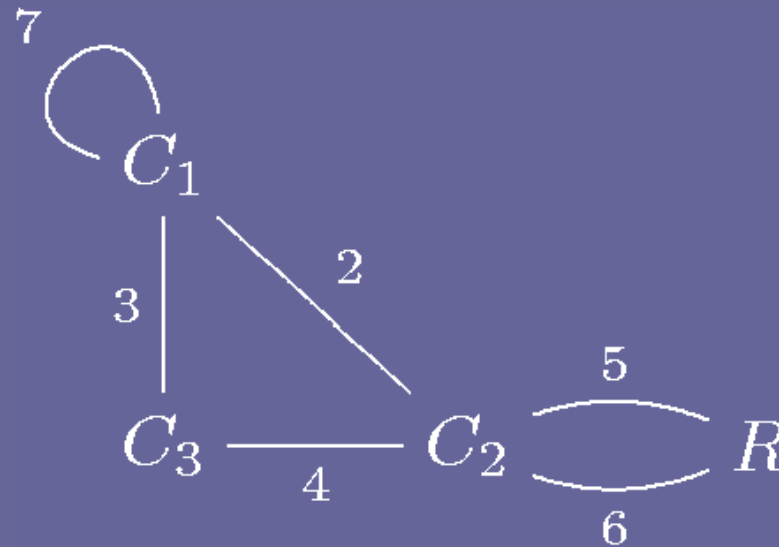
PointerComponentGraph



- A multigraph. Notice that loops and parallel edges are allowed.

PointerComponentGraph

- Which is:



- R is the linear connected component, and the others are cyclic components.

Result

For each legal string u ,

there is a successful reduction of u which applies snr on the set of pointers D
iff

D induces a spanning tree in the PC graph of u .

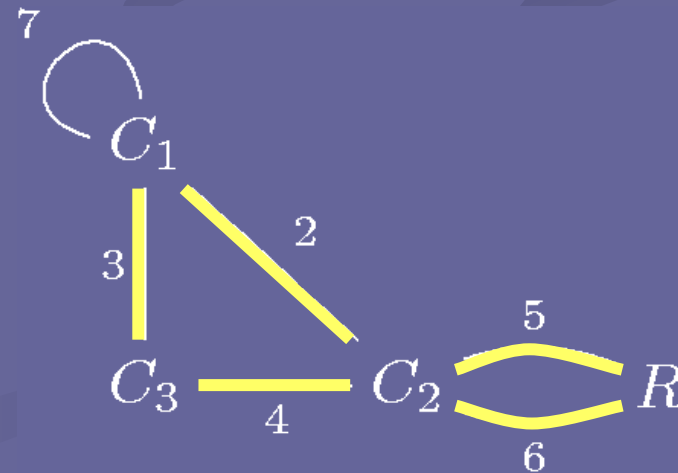
In the example: e.g.

$\{2,3,5\}, \{2,4,6\},$

but not

$\{2,5,6\}$ or $\{2,3\}$.

Also, 7 will never occur in such a set.

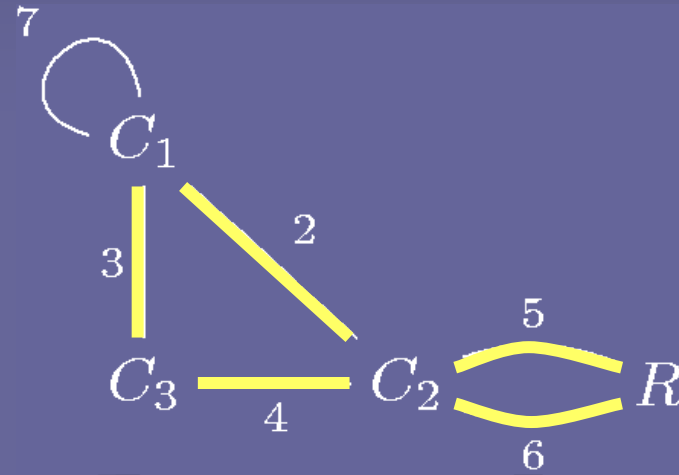


Example

$$u = 54372562\bar{7}346.$$

$D = \{2,3,5\}$ then, e.g.,
 $\varphi = snr_5 snr_2 snr_3 spr_7 sdr_{4,6}$ is a
successful reduction of u .

$D = \{2,4,6\}$ then, e.g.,
 $\varphi = snr_6 snr_2 spr_7 snr_4 sdr_{5,3}$ is a
successful reduction of u .



However, if we take $D = \{2,3,4\}$ then we always get stuck.

E.g., $(spr_5 spr_7)(u) = 62\bar{3}\bar{4}\bar{2}346$ No operation can be applied
on pointer 6 and no snr rule is applicable.

Nextquestion

We know:

- *How many* Snr rules are needed in successful reductions.
- *On which pointers* Snr rules are applied in successful reductions.

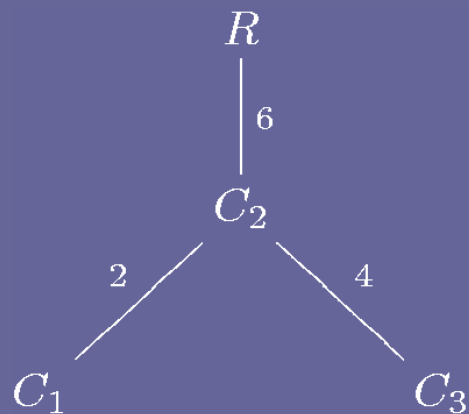
Next:

- What is *the order* of the Snr rules in successful reductions?

Extended result

If D induces a spanning tree T in the PC-graph, then the order in which snr rules are applied is determined by T with root R (the linear component).

Spanning tree in example :



Linear orderings



$(2,4,6) \rightarrow$ e.g. $\varphi = snr_6 snr_4 snr_2 spr_3 spr_5 spr_7$

$(4,2,6) \rightarrow$ e.g. $\varphi = snr_6 snr_2 spr_7 snr_4 sdr_{5,3}$

Reveals (in)dependence of string negative rules

Conclusion

Concept of pointer-component graph proves useful within the theory of gene assembly.

In biological terms: it allows for a characterization of applicability of loop recombination operations for a given micronuclear gene.

These characterizations correspond to efficient algorithms (making the graph and finding spanning trees in a graph are both computationally easy).

TheEnd

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Definitions

- Domain $dom(34\overline{3}\overline{2}) = \{2,3,4\}$.
- A composition of reduction rules φ is called a *reduction of u* if φ is applicable to u .

Example $\varphi = (spr_4 spr_3)$ is a reduction of $u = 342\overline{3}\overline{2}4$, since $\varphi(u) = \overline{2}2$.

- A reduction φ of u is called *successful* if $\varphi(u) = \lambda$.

Example $\varphi = (spr_2 spr_4 spr_3)$ is a successful reduction of $u = 342\overline{3}\overline{2}4$.