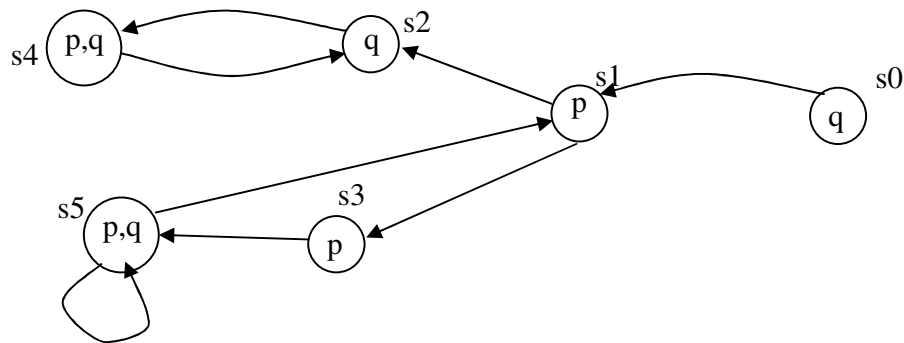


1. Alice and Bob are two persons seated at a table with one single fork. Each of them can be hungry and each of them can hold a fork. Write a CTL specification for the following properties:
 - a. Initially neither one of them holds the fork. [2 points]
 - b. They cannot hold the fork at the same time. [2 points]
 - c. If one of them gets hungry then that person will eventually pick up the fork. [2 points]
 - d. Each person cannot hold the fork continuously. [2 points]
 - e. Each person will be hungry infinitely often. [2 points]
 - f. Draw a transition system satisfying the above properties. [10 points]

2. Let p and q be two atomic propositions. Using the fixed-point method give the set of all states of the following transition system satisfying the CTL formula $EF(p \wedge AX q)$. [15 points]



3. Give an informal definition of operational, denotational and axiomatic semantics explaining their differences. [10 points]
4. Consider the command `iterate N times C end`, where N is a natural number and C is a command intended to be executed N times.
 - a. Formulate rules for the derivation of the execution

$$\langle \text{iterate } N \text{ times } C \text{ end}, \sigma \rangle \rightarrow \sigma'$$

in an initial state σ and terminating in the final state σ' . [10 points]

- b. Formulate a deduction rule for the partial correctness of the Hoare triple

$$\{ \phi \} \text{iterate } N \text{ times } C \text{ end } \{ \psi \}.$$

Give an informal justification of soundness. [10 points]

5. Calculate the weakest precondition of the following assignments with respect to the postcondition $a[i]=i$, where a is an array of positive integers, and i and j are two positive integers:

- a. $a[i] := a[i] + 1$ [5 points]
- b. $a[a[j]] := a[i] + 1$ [10 points]
- c. $a[a[i]] := a[j] + 1$ [10 points]

6. The following is a division algorithm. It assumes that x is non-negative and y is a positive integer. When the algorithm terminates, `quot` is the integer quotient and `rest` the remainder of x divided by y . Give a proof outline for the total correctness of the following Hoare triple:

```

{ x ≥ 0 ∧ y > 0 }
quot := 0;
rest := x;
while rest ≥ y do
    rest := rest - y;
    quot := quot + 1
od
{ ∃ n. 0 ≤ n < y ∧ quot * y + n = x }

```

[20 points]

The final score is given by the sum of the points obtained divided by 10 (with a maximum of 10).

Proof system for partial correctness:

- | | |
|---|--|
| <p>1. $\{ \phi \} \text{ skip } \{ \phi \}$</p> <p>3. $\frac{\{ \phi \} c1 \{ \phi \} \quad \{ \phi \} c2 \{ \psi \}}{\{ \phi \} c1 ; c2 \{ \psi \}}$</p> <p>5. $\frac{\{ \phi \wedge b \} c \{ \phi \}}{\{ \phi \} \text{ while } b \text{ do } c \text{ od } \{ \phi \wedge \neg b \}}$ </p> | <p>2. $\{ \phi[e/x] \} x := e \{ \phi \}$</p> <p>4. $\frac{\{ \phi \wedge b \} c1 \{ \psi \} \quad \{ \phi \wedge \neg b \} c2 \{ \psi \}}{\{ \phi \} \text{ if } b \text{ then } c1 \text{ else } c2 \text{ fi } \{ \psi \}}$</p> <p>6. $\frac{\phi \Rightarrow \phi1 \quad \{ \phi1 \} c \{ \psi1 \} \quad \psi1 \Rightarrow \psi}{\{ \phi \} c \{ \psi \}}$ </p> |
|---|--|