

**Question 1:** [1,5 points]

- Find at least four different pairs of languages  $A$  and  $B$  over the alphabet  $\Sigma = \{0, 1\}$  for which their concatenation  $AB$  is the set of strings  $\{10, 111, 1010, 1000, 10111, 101000\}$ .
- Give regular expressions for the intersection, union, and concatenation of the two languages  $A = \{w \in \Sigma^* \mid w \text{ begins with } 11\}$  and  $B = \{w \in \Sigma^* \mid w \text{ ends with } 00\}$ .
- Let  $R$  be a regular expression. Do  $R + \emptyset$  and  $R\Lambda$  denote the same language? And what about  $R\emptyset$  and  $\emptyset R$ ? And  $\Lambda + R(R^*)$  and  $R^*$ ? Justify your answers.

**Question 2:** [2,5 points]

- Give a non-deterministic finite automaton with at most three states accepting the language of the regular expression  $\Lambda + ((0 + 00^*1)1^*)$ .
- Using the powerset construction, find a deterministic finite automaton that recognizes the same language as the non-deterministic finite automaton in the above item a).
- Suppose  $M = (Q, \Sigma, q_0, A, \delta)$  is a non-deterministic finite automaton with  $\Lambda$ -transitions recognizing the language  $L$ . Let  $M_1$  be the automaton obtained from  $M$  by adding a  $\Lambda$ -transition from each state in  $A$  to  $q_0$ . Describe the language  $L(M_1)$  in terms of  $L$ .

**Question 3:** [2,5 points]

For each statement below decide if it is true or false. If it is true, prove it, otherwise give a counterexample.

- If  $L_1 \subseteq L_2$  and  $L_1$  is not regular, then also  $L_2$  is not regular.
- If  $L$  is not regular then also its complement  $L'$  is not regular.
- If  $L_1$  is regular,  $L_2$  is not regular and  $L_1 \cap L_2$  is regular, then  $L_1 \cup L_2$  is not regular.

**Question 4:** [2 points]

- Is the language of all odd-length strings over the alphabet  $\{a,b,c\}$  regular? Justify your answer.
- Is the language of all odd-length strings over the alphabet  $\{a,b,c\}$  with middle symbol  $b$  regular? Justify informally your answer.
- Use the pumping lemma for regular languages to show that the language  $L = \{a^n b^m c^k \mid n+m < k\}$  is not regular.

**Question 5:** [1,5 points]

- Give the definitions of context free grammar and of regular grammar.
  - Find a *context-free* grammar for language  $L_1 = \{a^n b^m c^k \mid n+2m \geq k\}$ .
  - Find a *regular* grammar for the language  $L_2 = \{a^n b^m \mid n \geq 0, m \geq 2\}$ .
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The final score is given by the sum of the points obtained.