

TILES

from pattern to computation

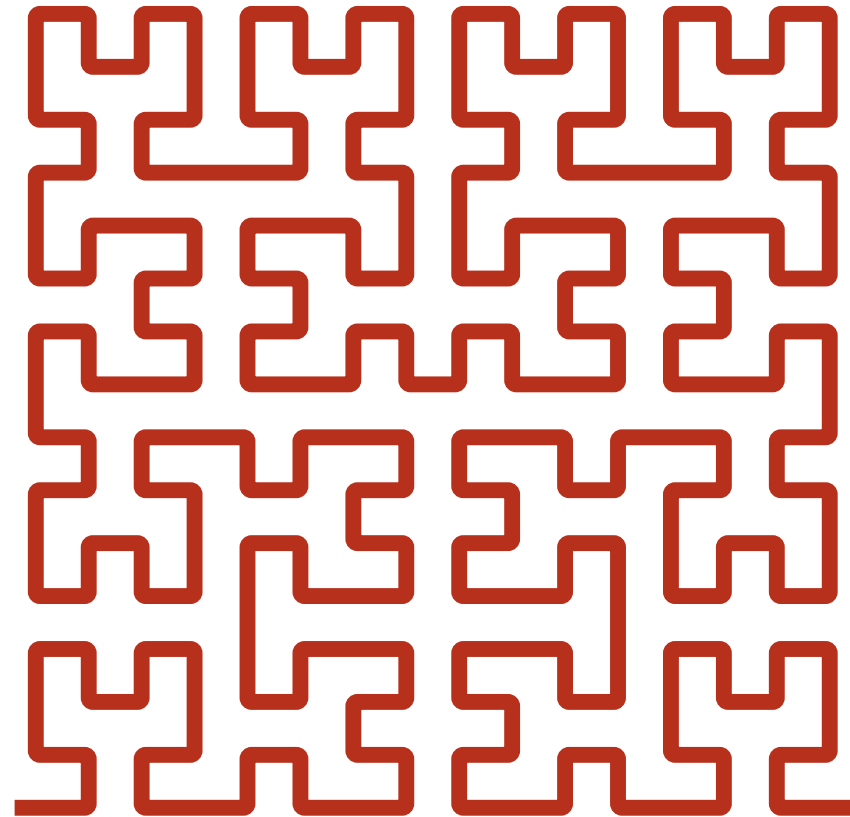
Hendrik Jan Hoogeboom

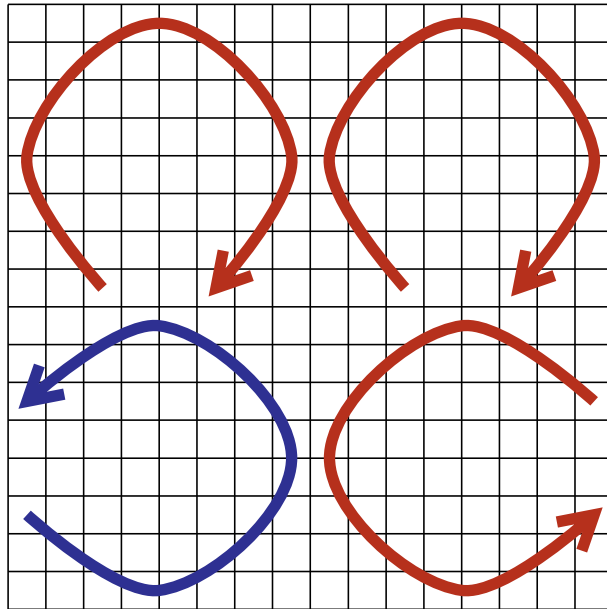
Universiteit Leiden, LIACS

www.liacs.nl/~hoogeboo/praatjes/tegels/

■ **personal motivation**

- 1-dim line
- 2-dim pattern



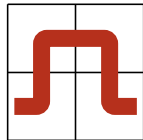
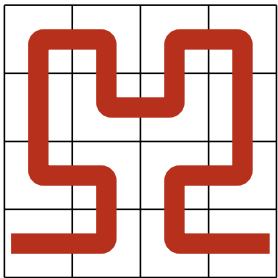
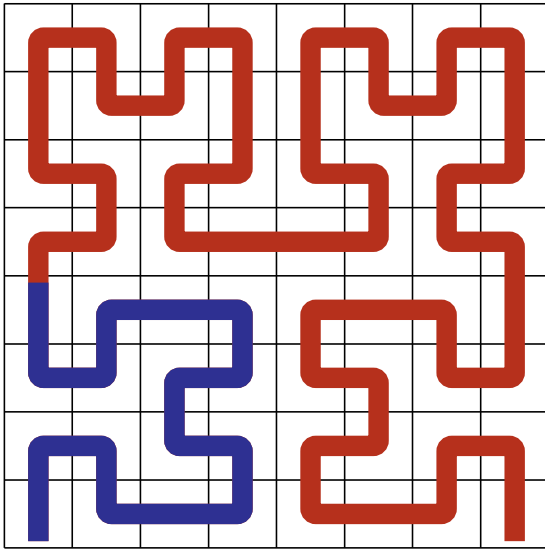


wikipedia

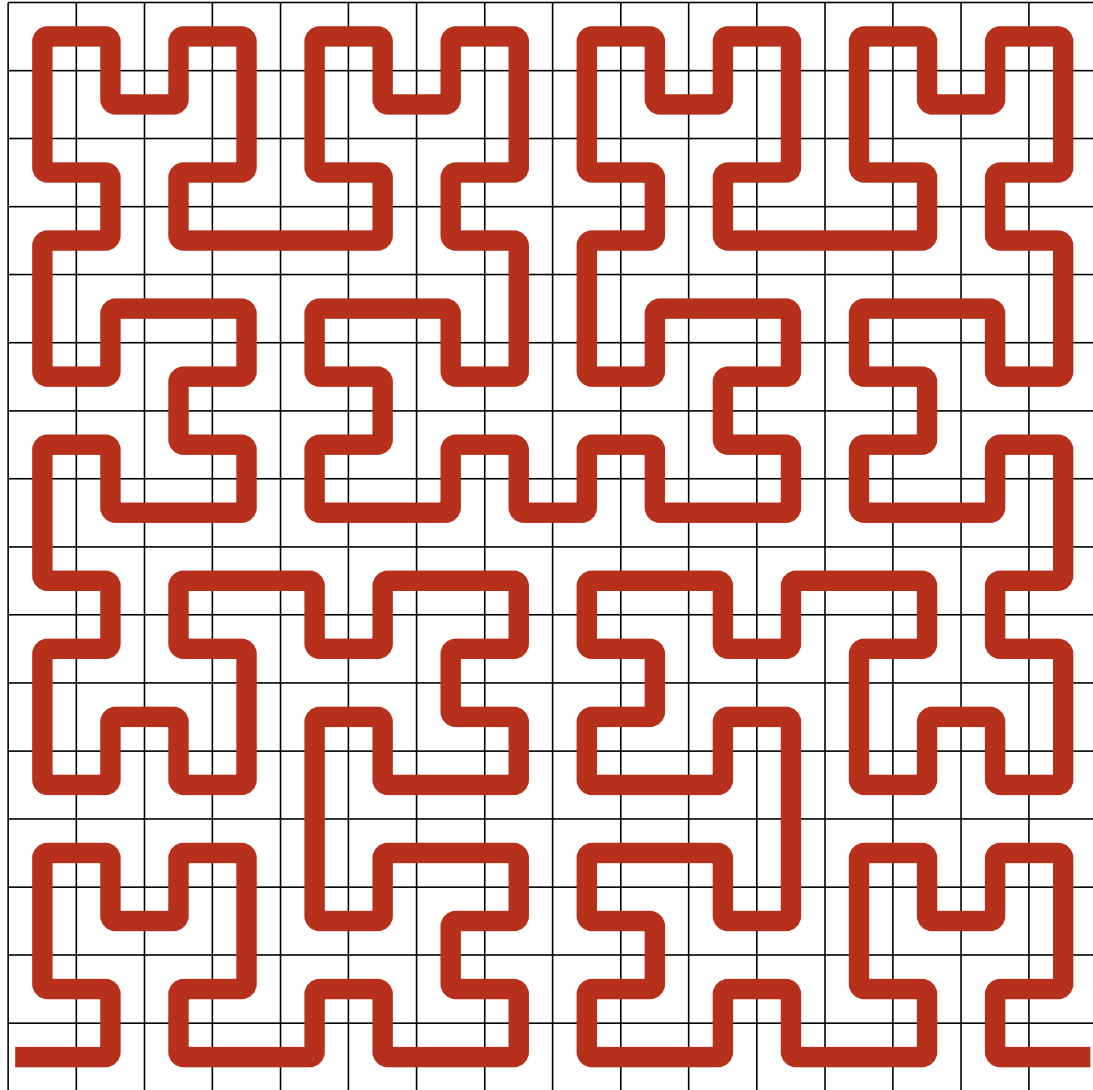
```
to starthilbert :size :level
  hilbert (:size / power 2 (:level-1)) :level 1
end
```

```
to hilbert :size :level :parity
  if :level = 0 [stop]
  right 90 * :parity
  hilbert :size (:level-1) (:parity * -1)
  forward :size
  right -90 * :parity
  hilbert :size (:level-1) :parity
  forward :size
  hilbert :size (:level-1) :parity
  right -90 * :parity
  forward :size
  hilbert :size (:level-1) (:parity * -1)
  right 90 * :parity
end
```

```
call starthilbert 200 5
```



six 'tiles':
two lines, four curves

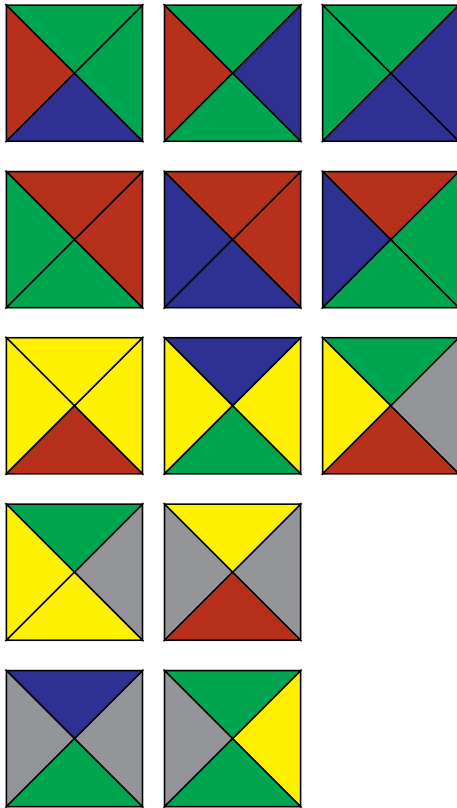


■ background

input: set of tiles

question: do they tile the plane?
(a rectangle)

matching edges, no rotations



aperiodic tiling pattern (hard)

Karel Culik II, 1996

1961 Wang decidable

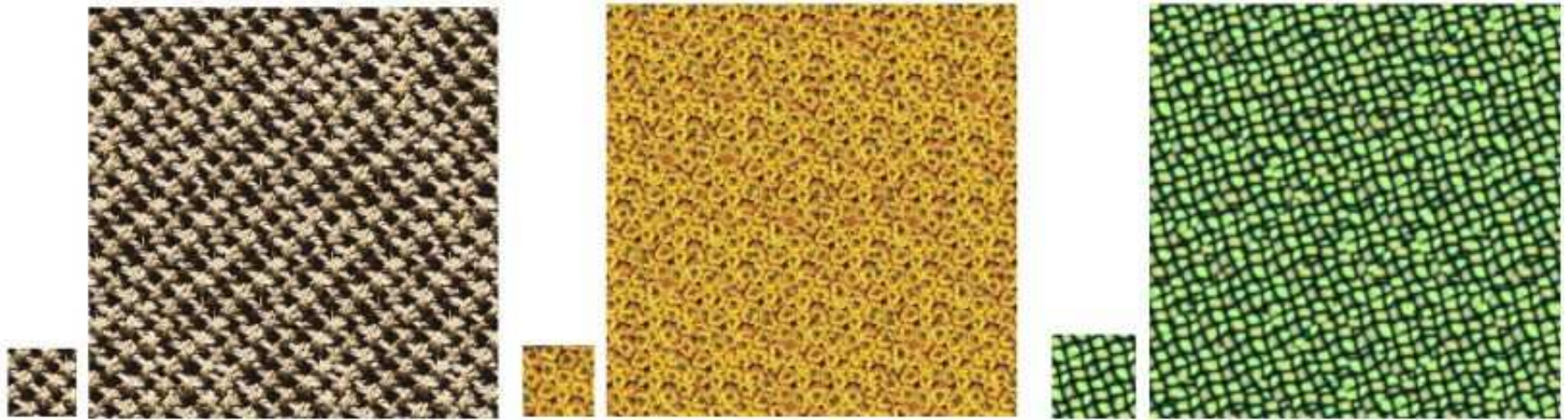
assumption: tiling is periodical

1966 Berger undecidable

aperiodic tiling (with 20,426 tiles)

1996 Culik 13 tiles

1974 Penrose 2 shapes (kite & dart)



Cohen, Shade, Hiller, and Deussen: Wang Tiles for Image and Texture Generation. ACM Transactions on Graphics 22 (2003) 287-294.

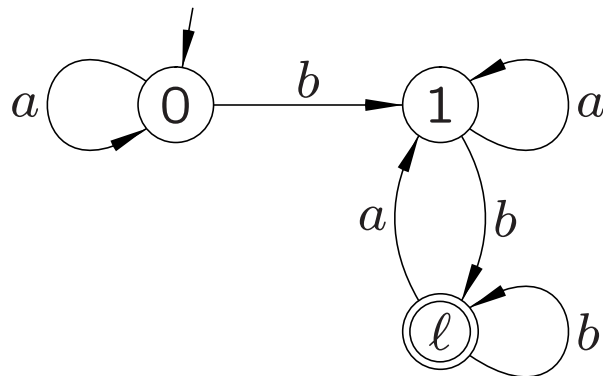
Ares Lagae: Wang Tiles in Computer Graphics, 2009

aim: picture specifications

two-dimensional theory

strings: Chomsky hierarchy

grammars, automata, logic, ...

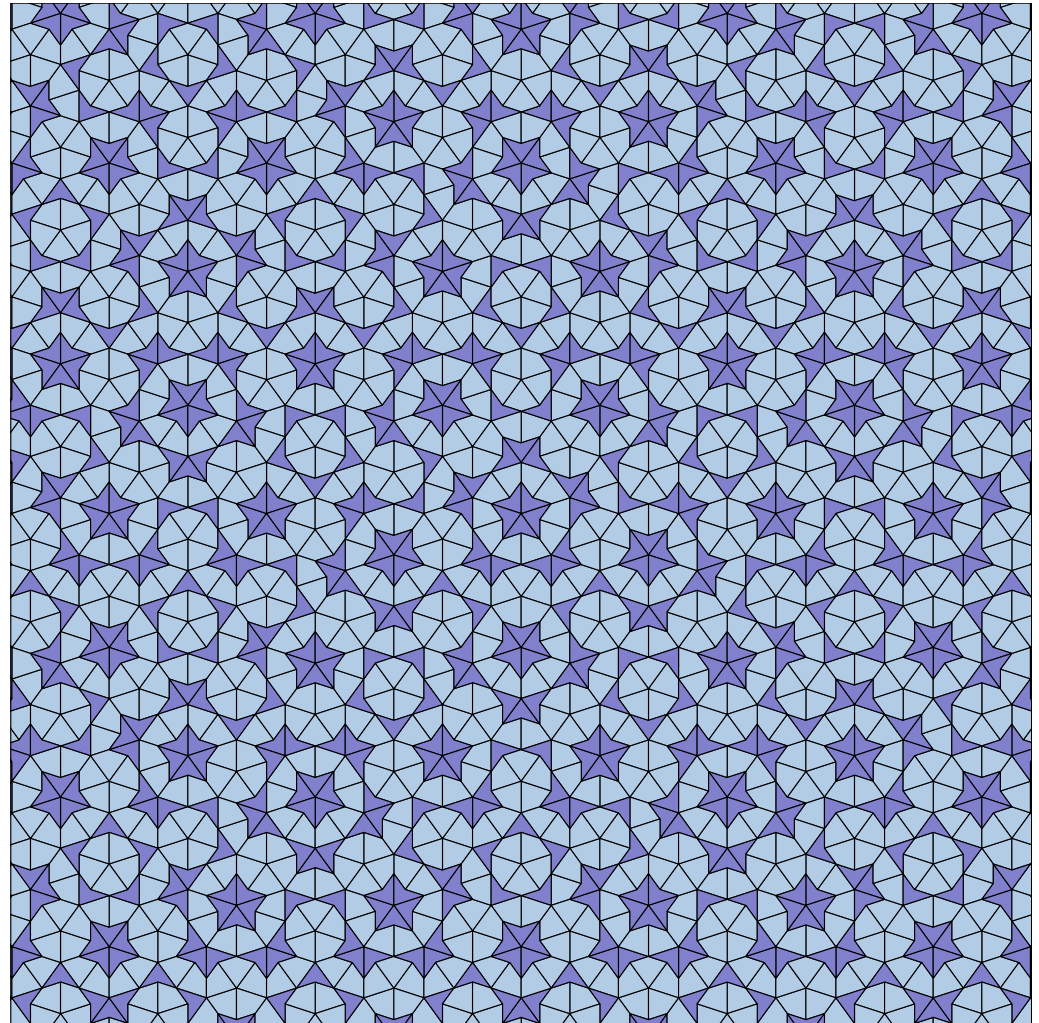


finite state specification

FSA

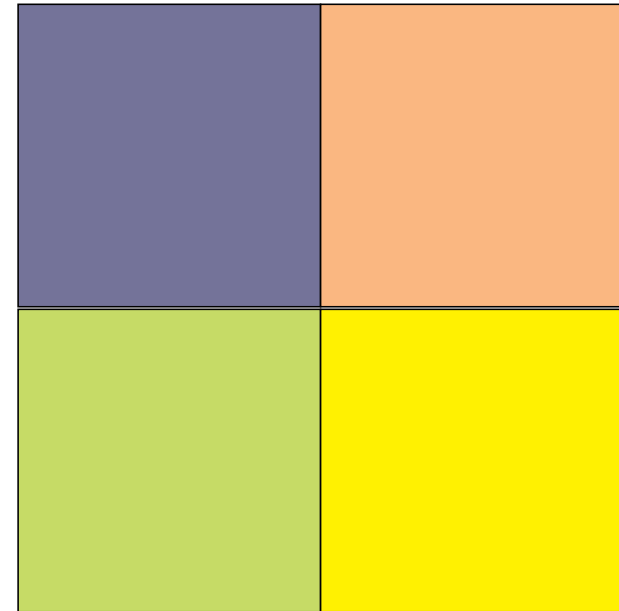
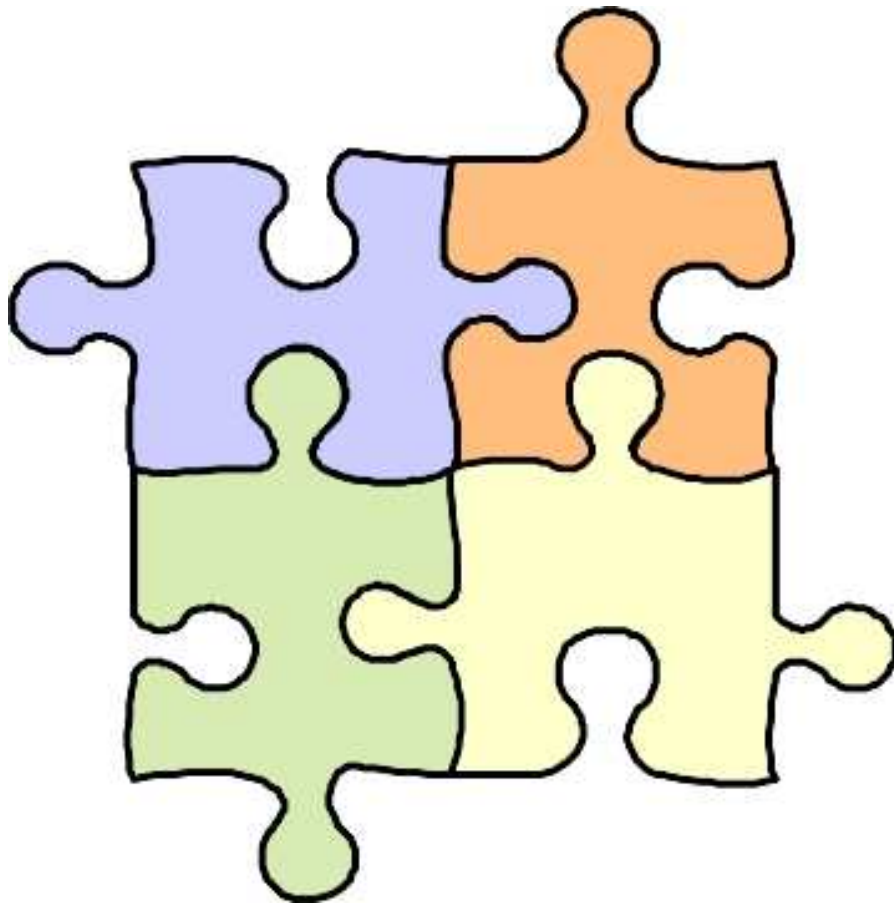
- (non-)determinism
- closure properties
 - union, intersection, iteration, complement, ...
- decision properties
 - emptiness, inclusion, ...
- regular languages
 - automata, expressions, grammars, logic...

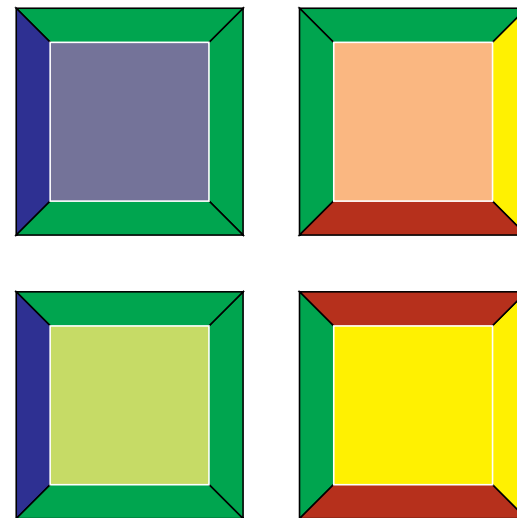
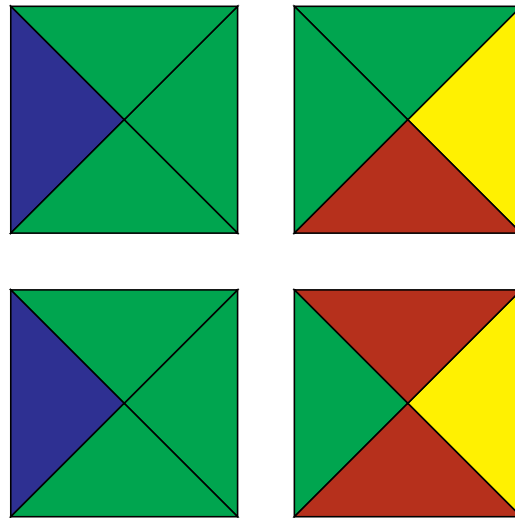
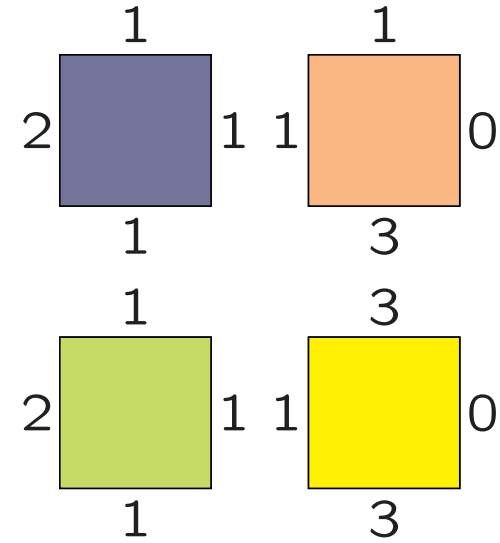
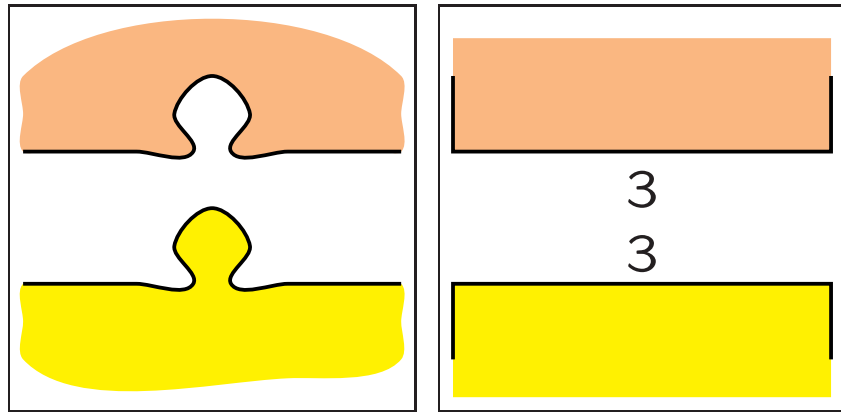
■ Rules of the game

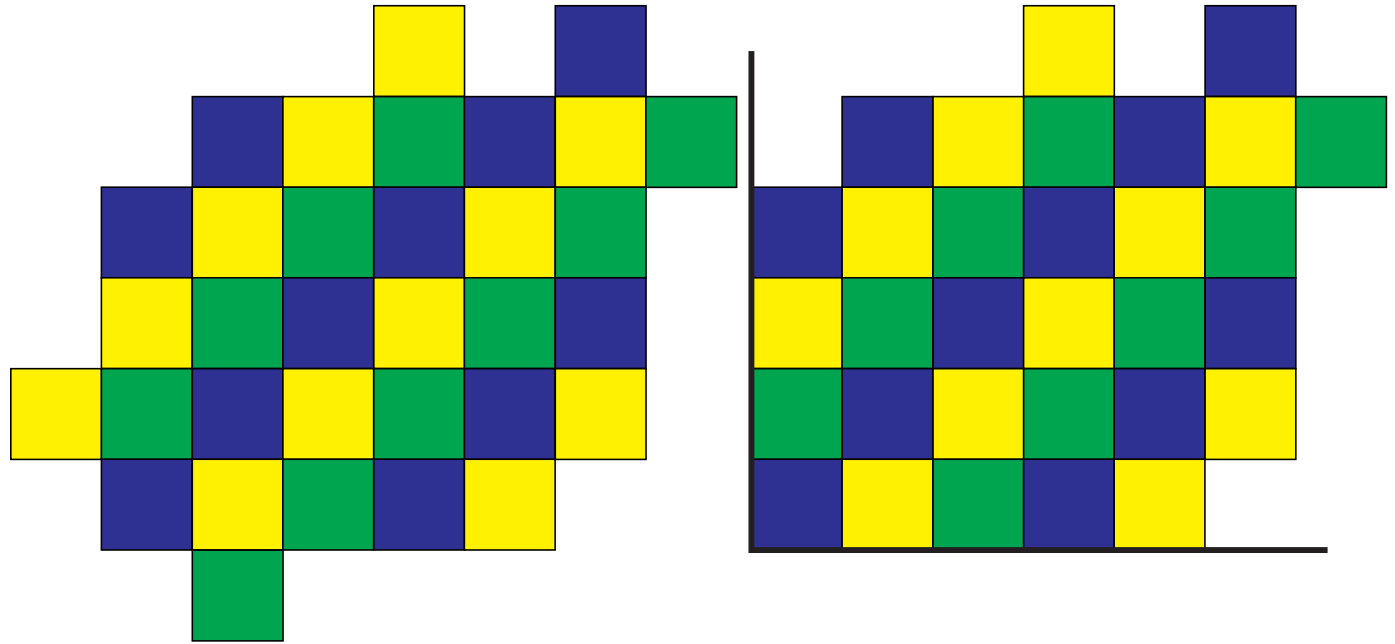
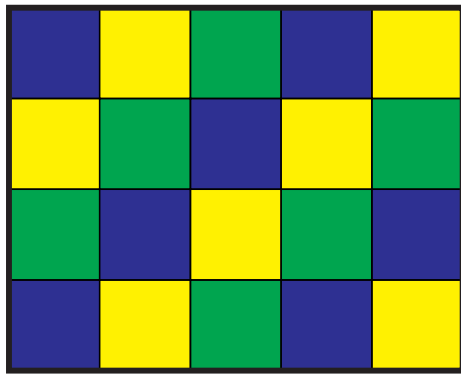


Penrose (1974)

© Franz Gähler, Stuttgart

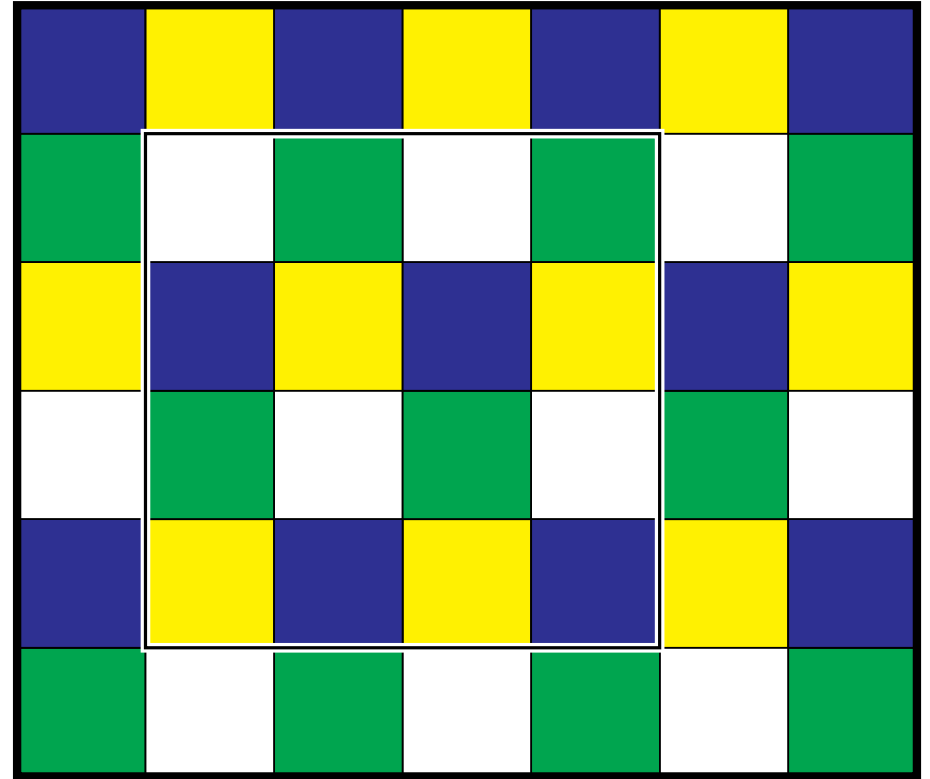
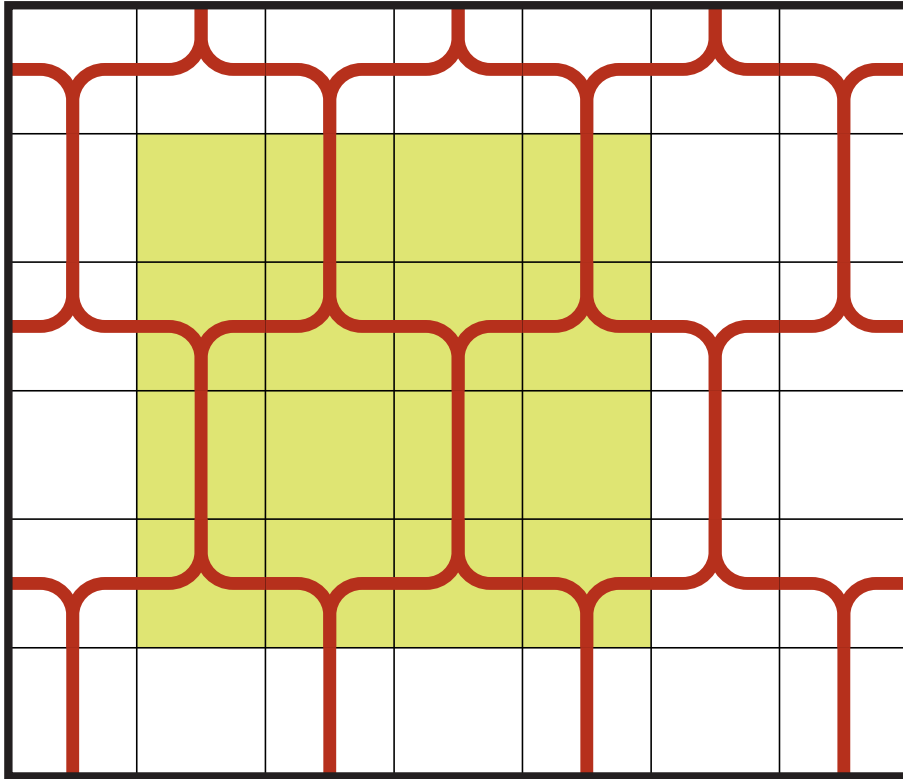




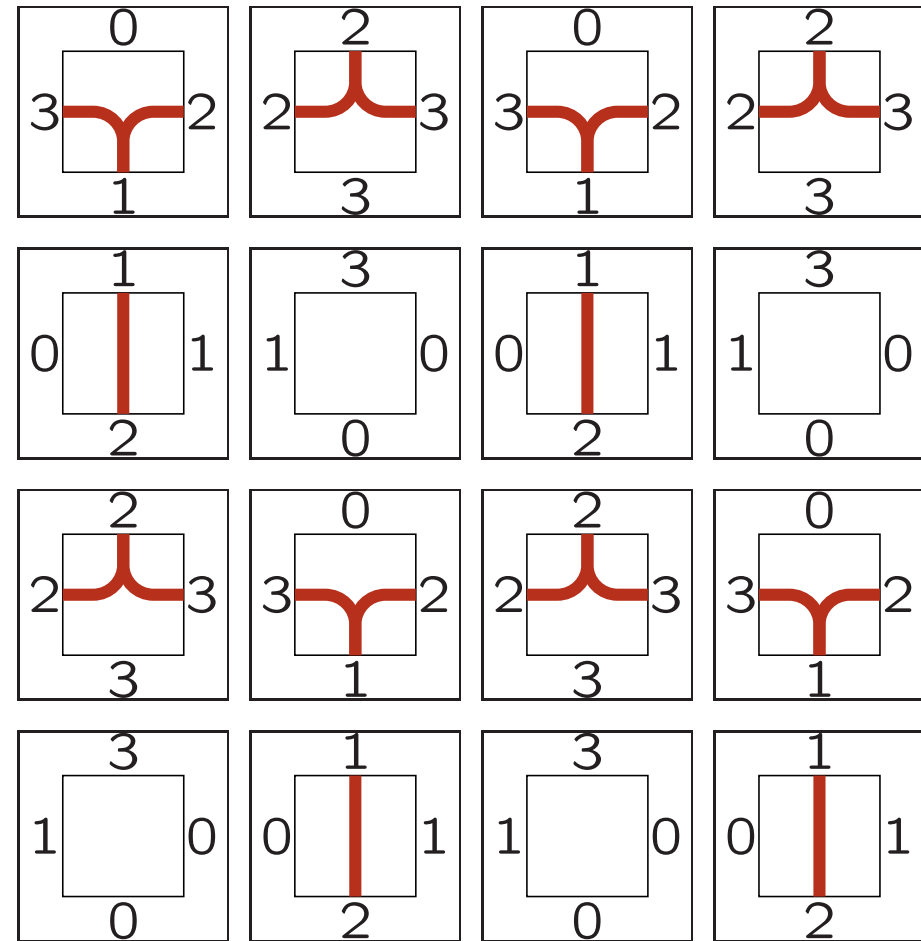
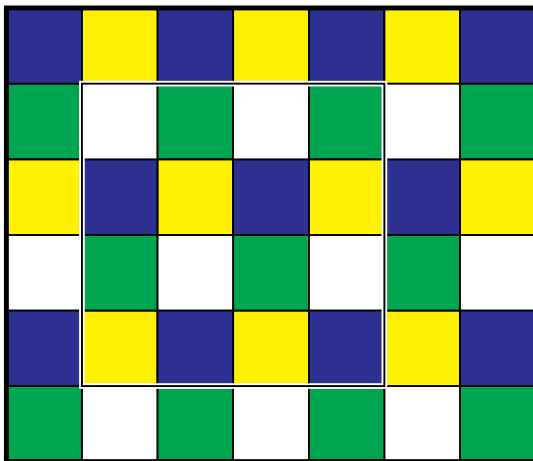
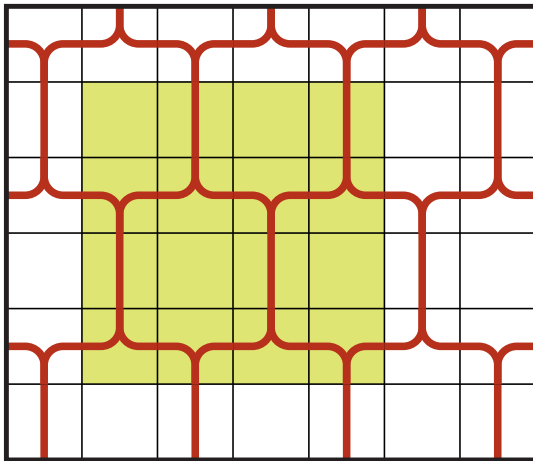


- rectangle
- (half) plane
- quadrant

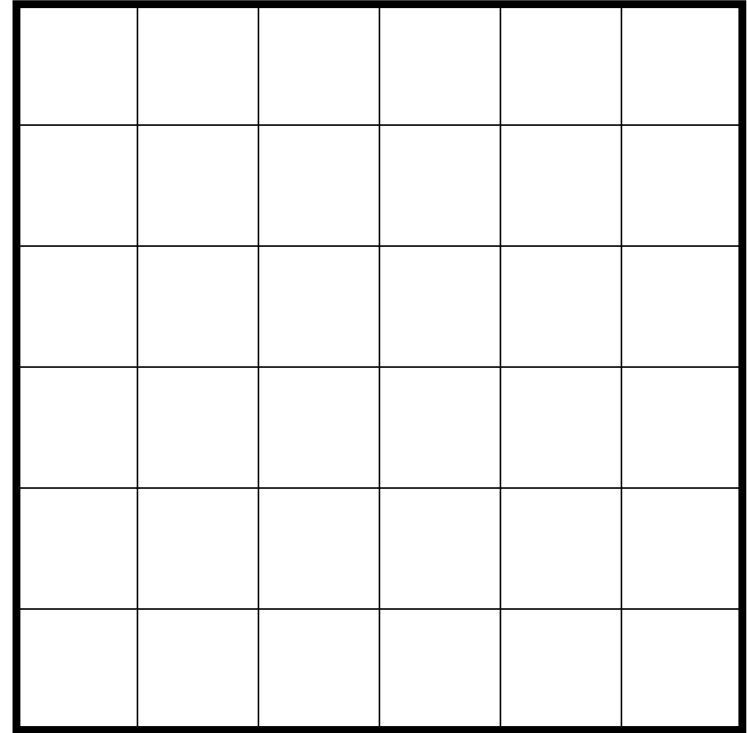
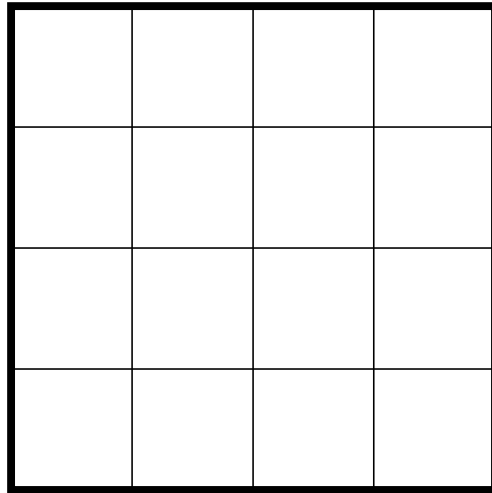
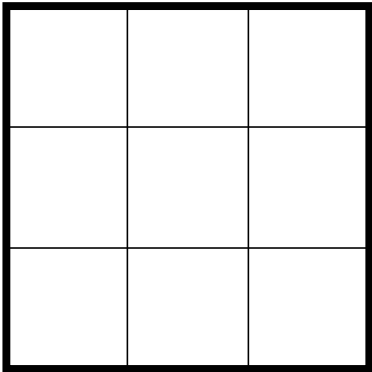
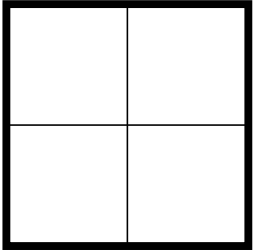
■ Examples



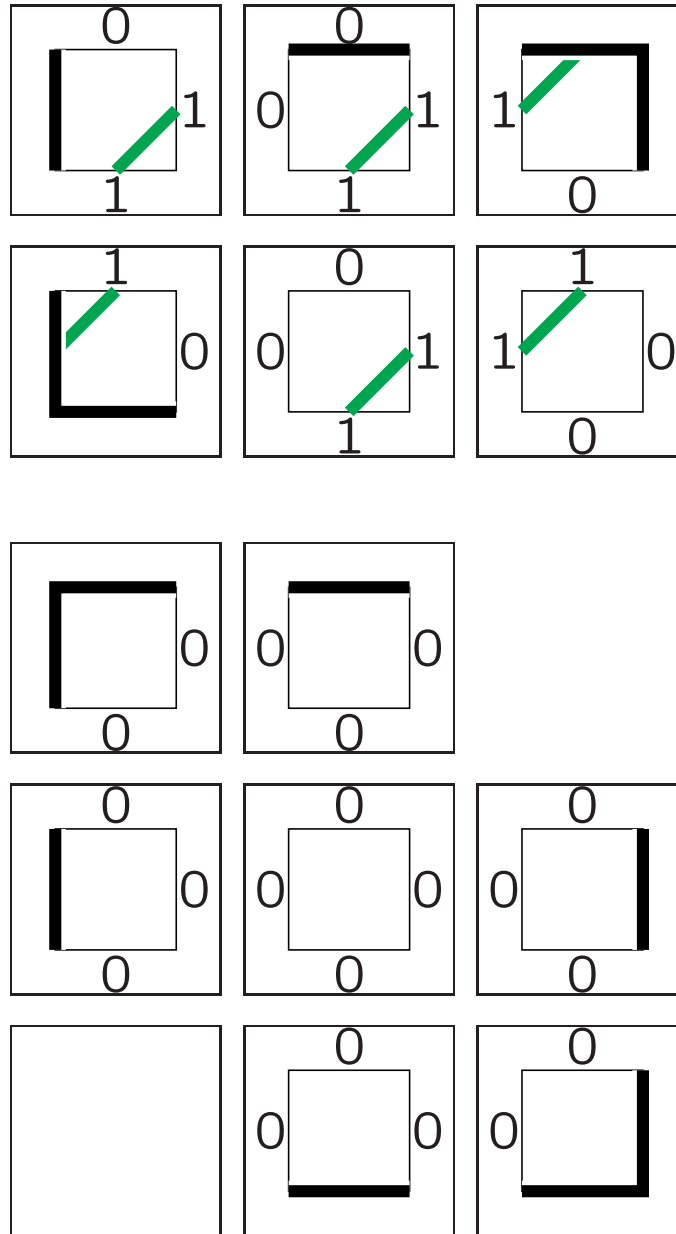
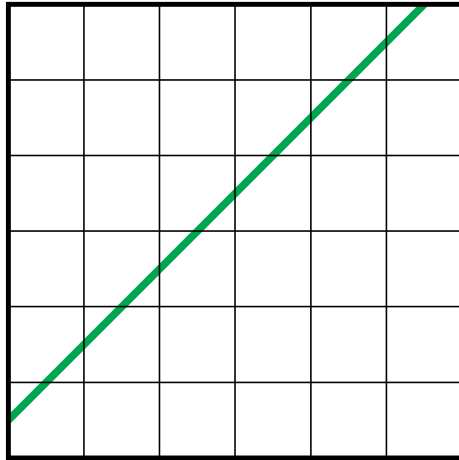
four 'colours'

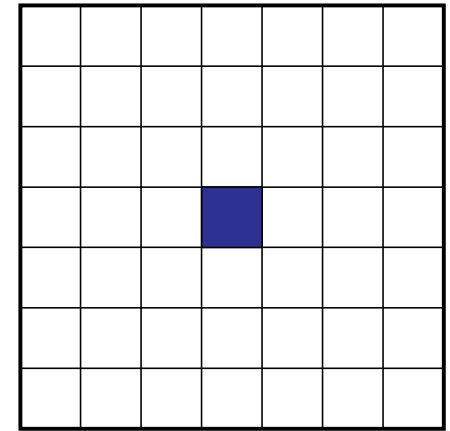
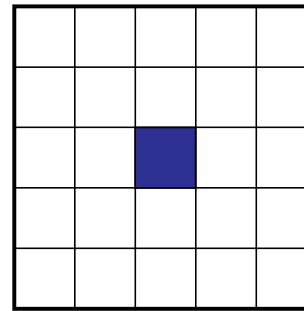
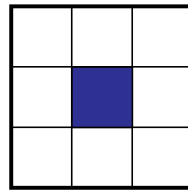
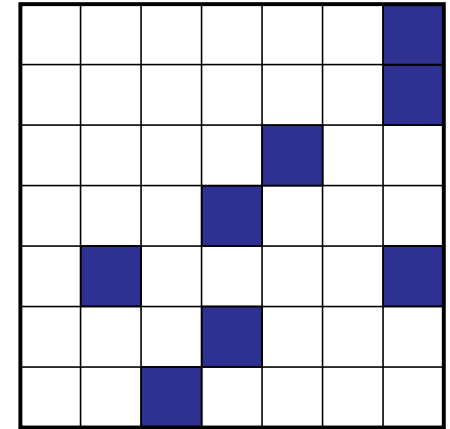
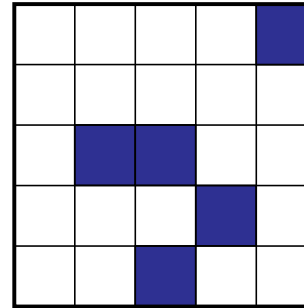
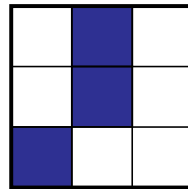


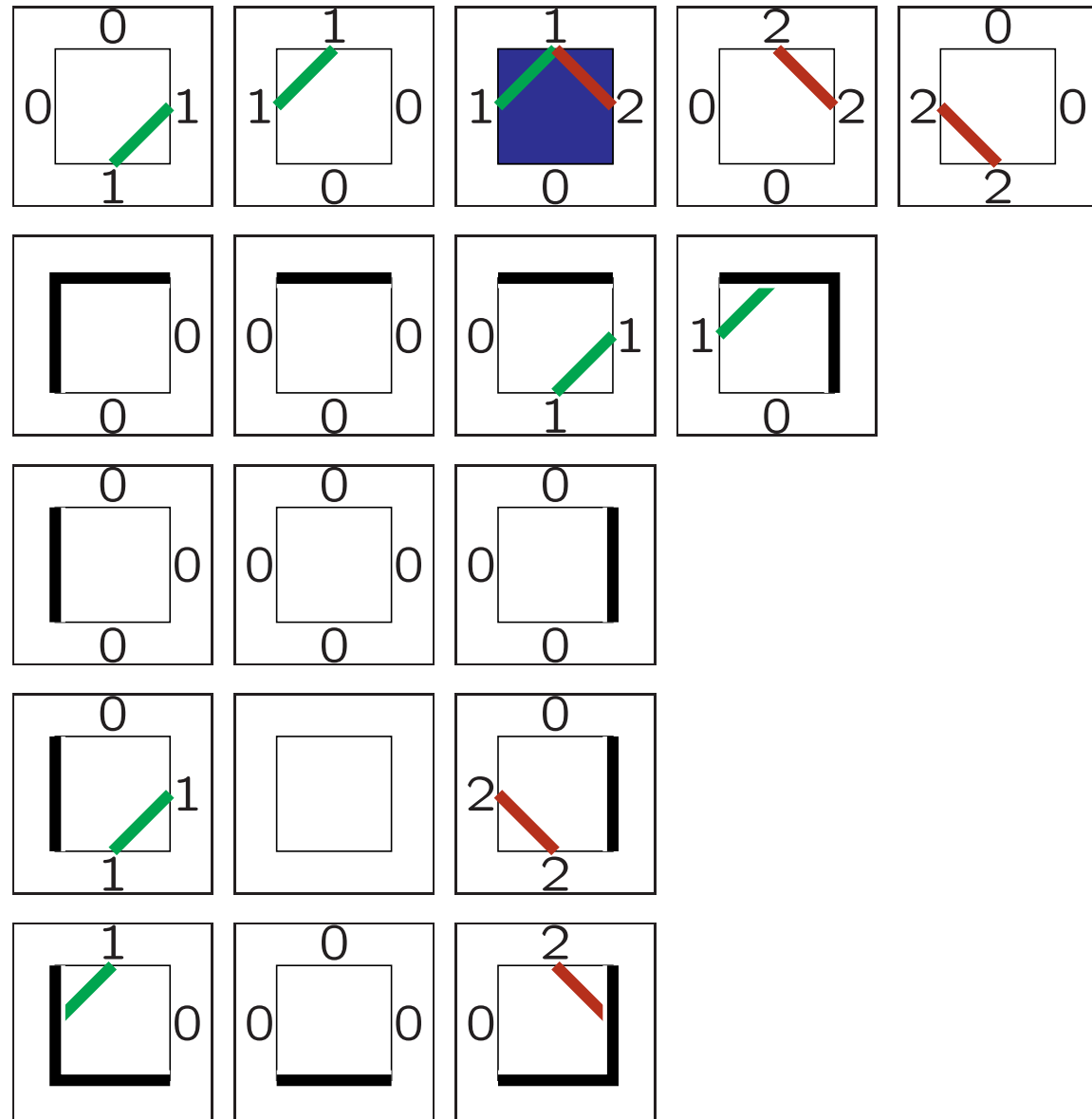
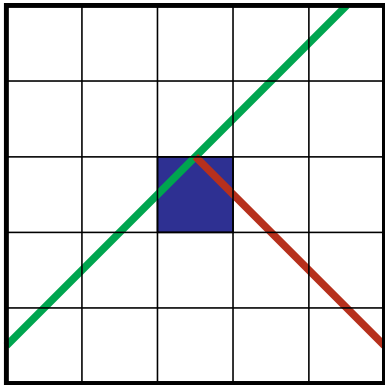
four tiles
(and some borders)

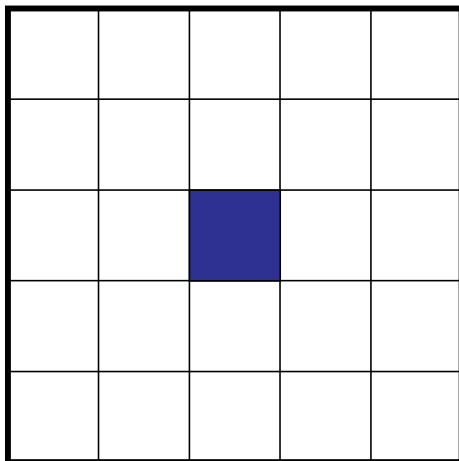
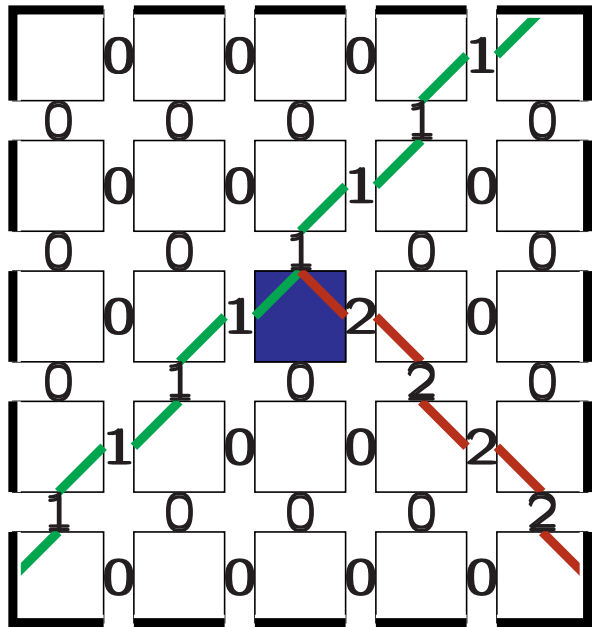


etcetera ...









tiling system:

$(\Sigma, \Gamma, T, c, \varphi)$

Σ, Γ tile and edge colours

T tiles with four-sided markings

$c \in \Gamma$ border marking

$\varphi : T \rightarrow \Sigma$ tile \mapsto colour

tiling: markings (in Γ) match

Wang picture language

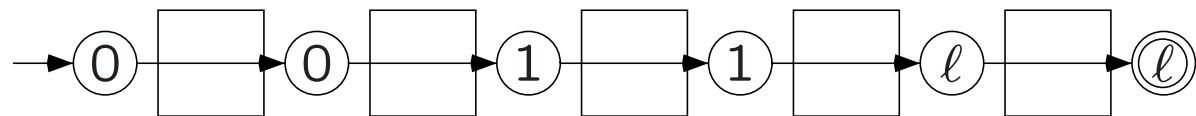
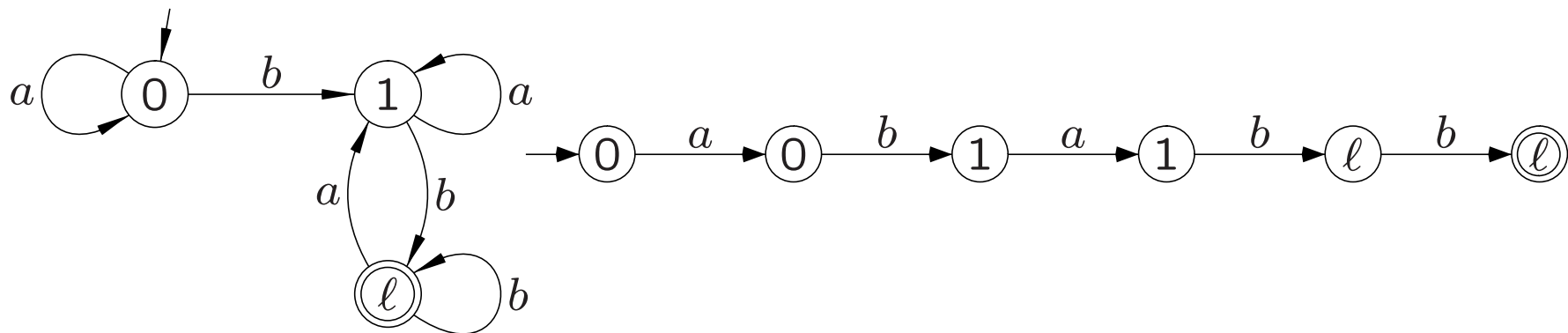
Giammarresi, Restivo:

REC – recognizable

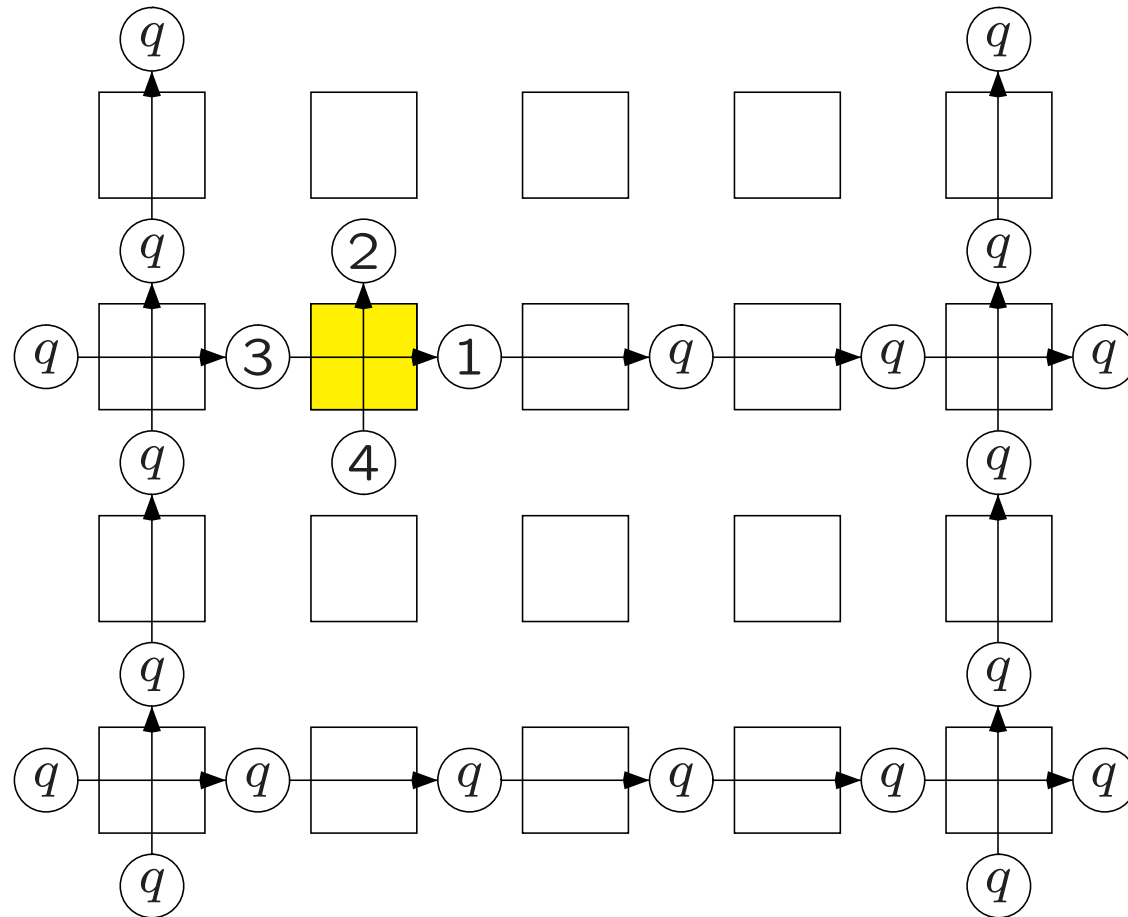
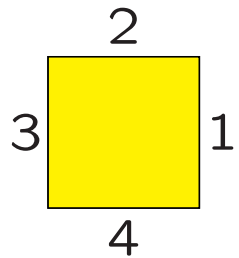
local lattice languages

h(LL)

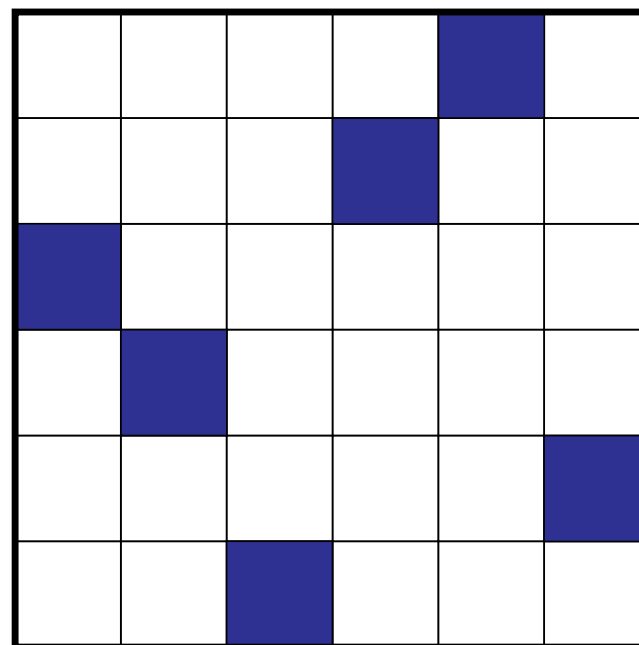
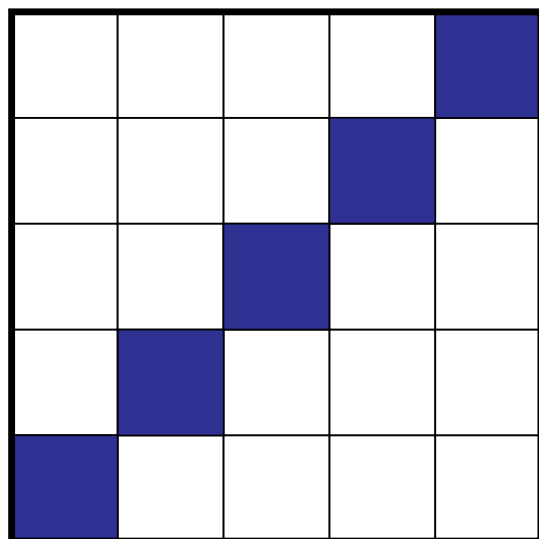
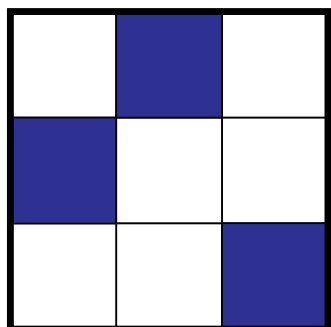
computation: finding transitions

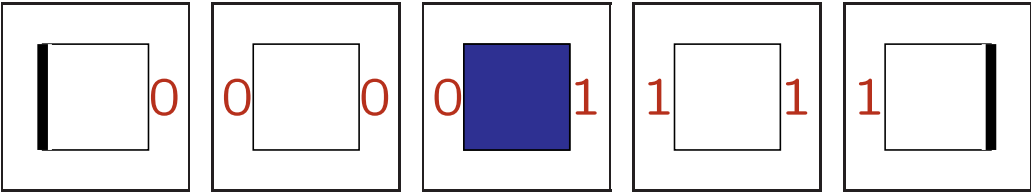
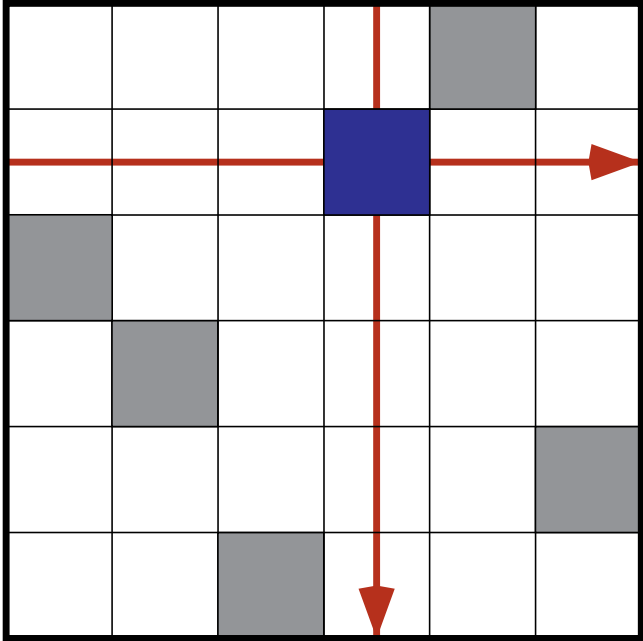


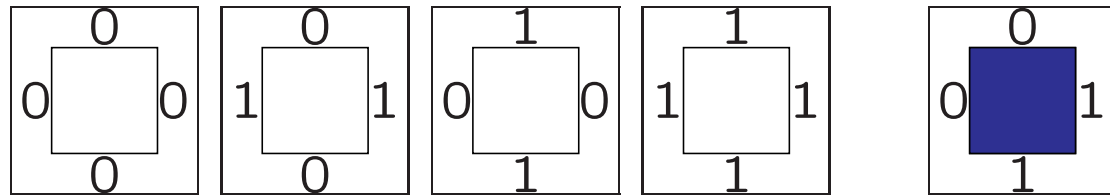
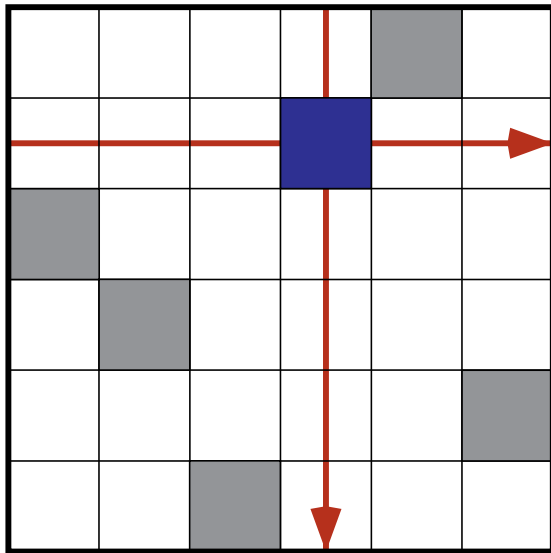
tile \sim transition



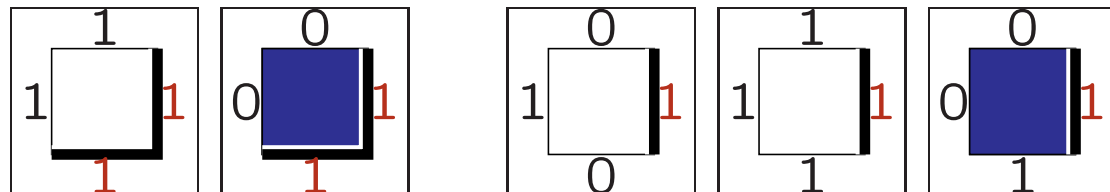
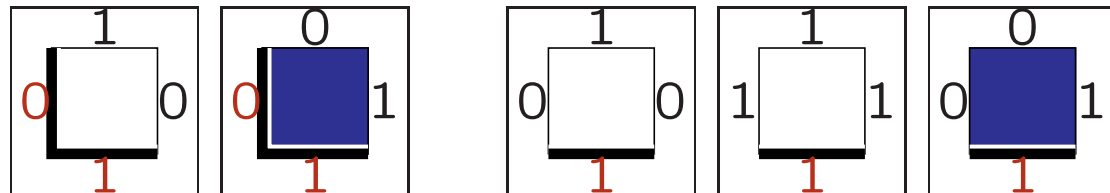
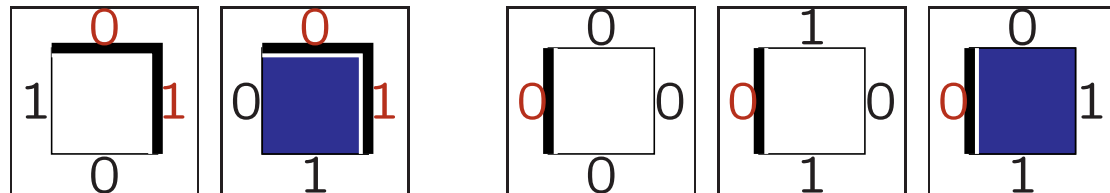
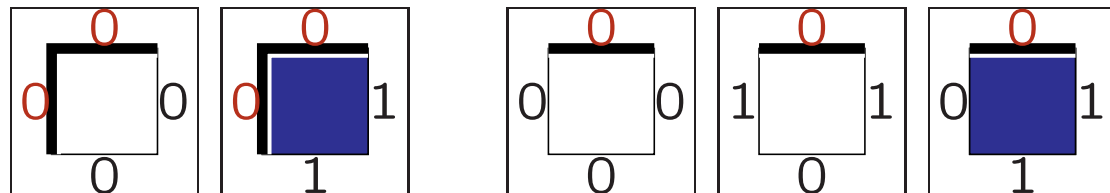
note: hor+vert transitions 'connected'







left/top 0 right/bottom 1

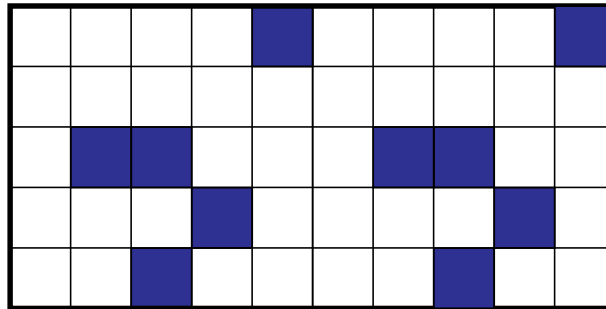


closed under

- renaming
- concatenation, iteration (hor & vert)
- union, intersection
- rotation
- *not* closed under complement

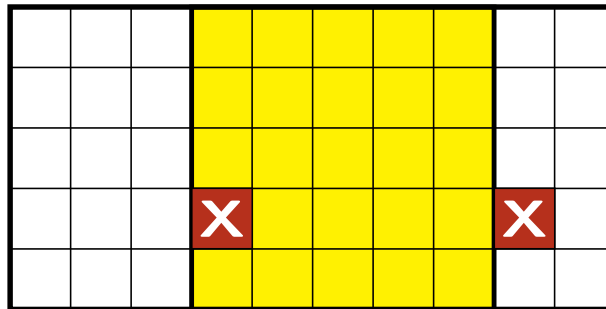
- *not* closed under complement

$\{ww \mid w \in \{a, b\}^*\}$
not context-free . . .

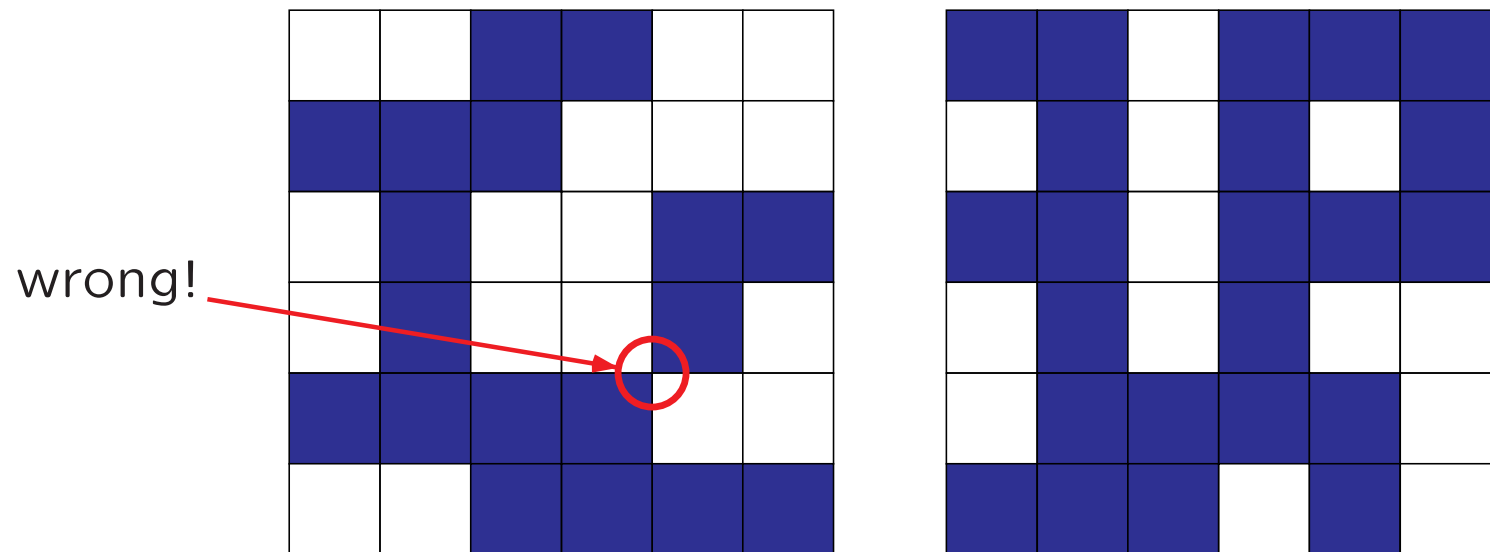


c tile colours, e edge colours
choose $c^{n^2} > e^n$, cut & glue

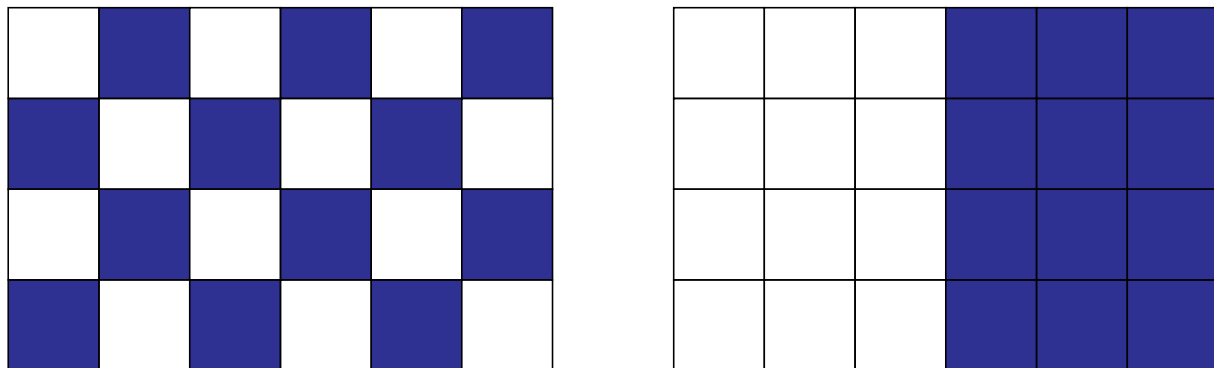
but complement is



- blue tiles are connected (*difficult*)

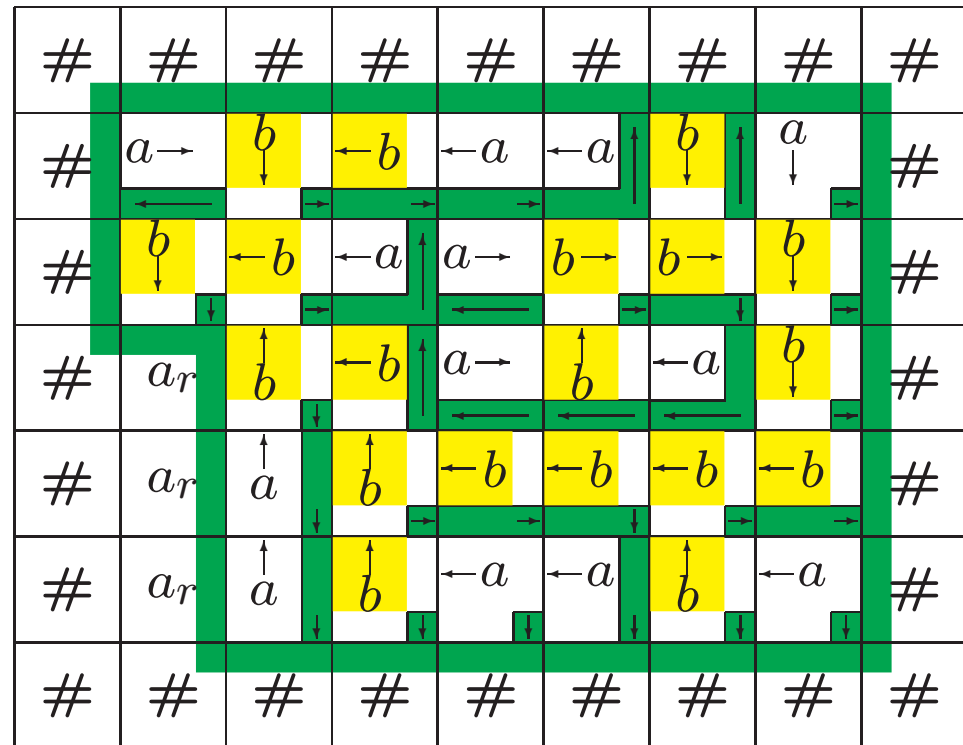


- equal numbers (*very difficult*)



	a	b	b	a	a	b	a	
	b	b	a	a	b	b	b	
	a	b	b	a	b	a	b	
	a	a	b	b	b	b	b	
	a	a	b	a	a	b	a	

using Wang tiles?



Klaus Reinhardt

On Some Recognizable Picture-Languages (1998)

logic **MSO** \Leftrightarrow finite state automata

sets (of positions)

$$\begin{array}{l}
 \text{even number} \\
 \text{of } a\text{'s} \\
 (\exists X) [\\
 \quad (\forall k)(k \in X \rightarrow a(k)) \\
 \quad \wedge (\forall i)(\text{first}_a(i) \rightarrow i \in X) \\
 \quad \wedge (\forall i)(\forall j)(\text{next}_a(i, j) \rightarrow (i \in X \leftrightarrow j \notin X)) \\
 \quad \wedge (\forall i)(\text{last}_a(i) \rightarrow i \notin X) \\
]
 \end{array}$$

where $\text{first}_a(i)$ means

$$a(i) \wedge \neg(\exists j)(a(j) \wedge j < i)$$

etc.

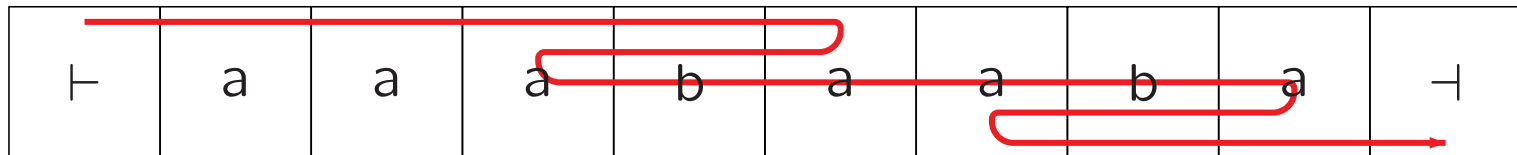
b (a) a b (a) b b a b (a) a b

Prop. **EMSO** = **WANG**

■ **four-way automata**

Prop. **two-way fsa** are equivalent to one-way aut.
(deterministic and non-deterministic)

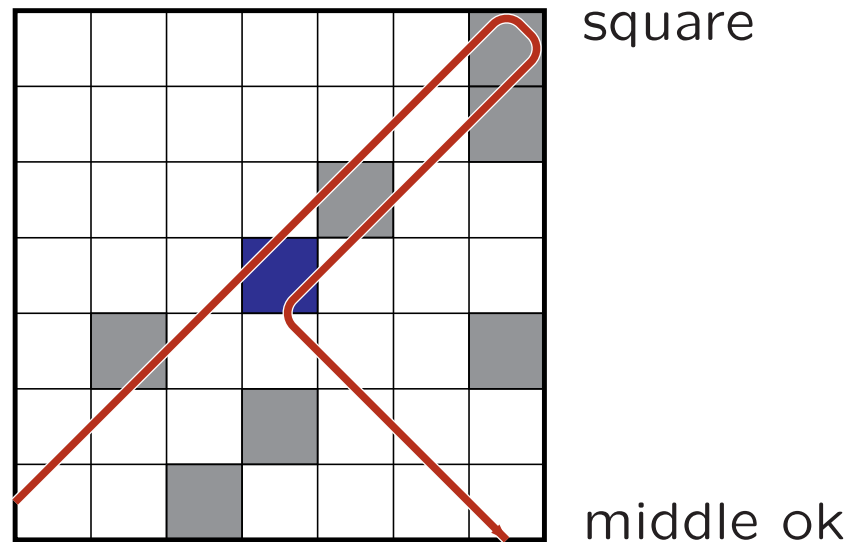
Rabin&Scott, Shepherdson (1959)



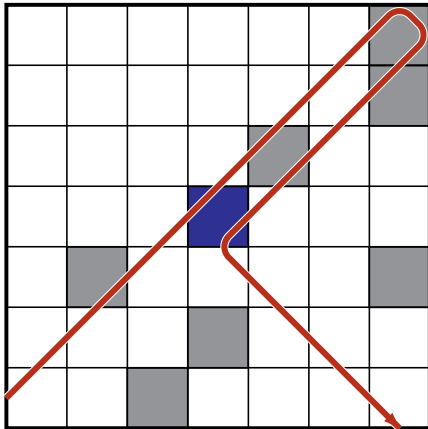
Blum, Hewitt: Automata on a 2-dimensional tape
(1967)

can a 'robot' recognize a tiling?

state \times colour \longrightarrow state \times direction



non-deterministic – guess middle



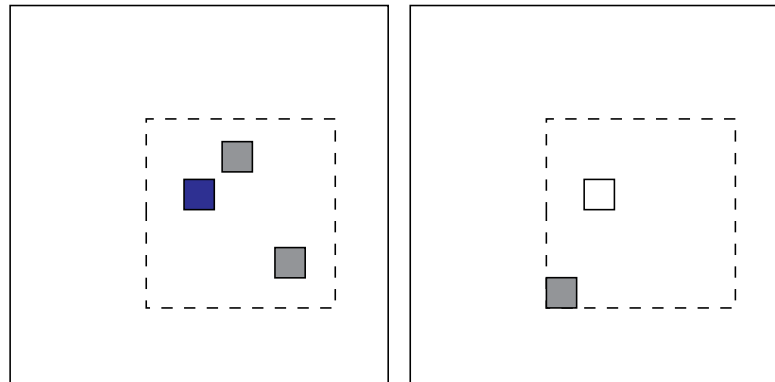
only *non*-deterministically

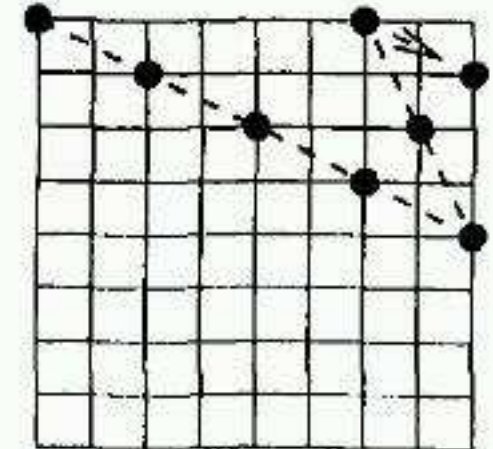
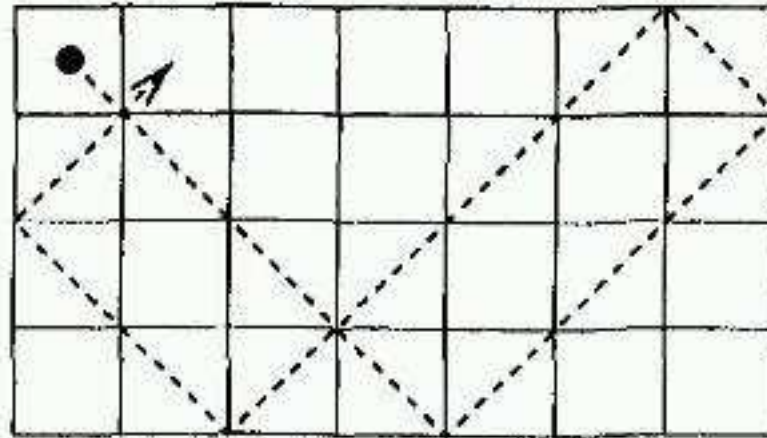
s states, c colours

in-state, out-state (elsewhere)

choose $(4m \cdot s)^{4m \cdot s} < 2^{m^2}$

two equivalent squares



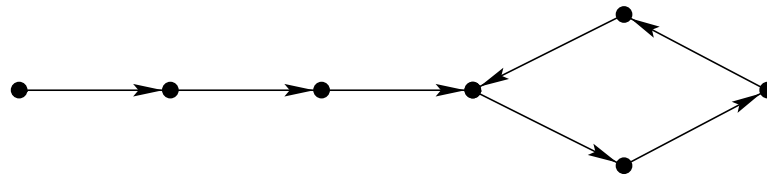


Lindgren, Moore, and Nordahl

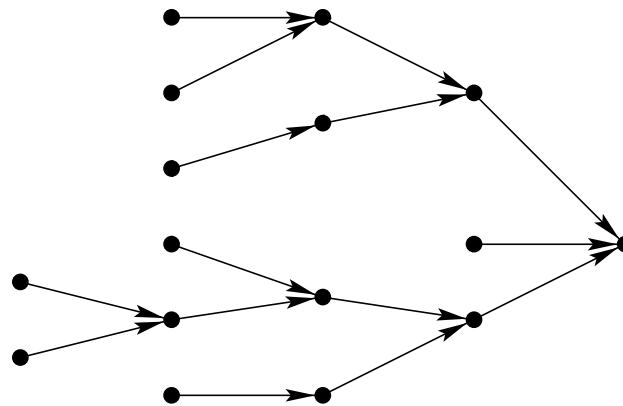
"By bouncing like a billiard ball or making knights' moves, and ending one cell from the corner, a DFA can check that the two sides of a rectangle are **mutually prime**, or that the side of a square is a **power of 2**."

"bouncing like a billiard ball ... can check that ..."

⇒ avoid loops!



simulate search space **backwards**: a tree!



this can be done by a 4DFA !

Sipser, Halting space-bounded computations (1980)

deterministic closed under

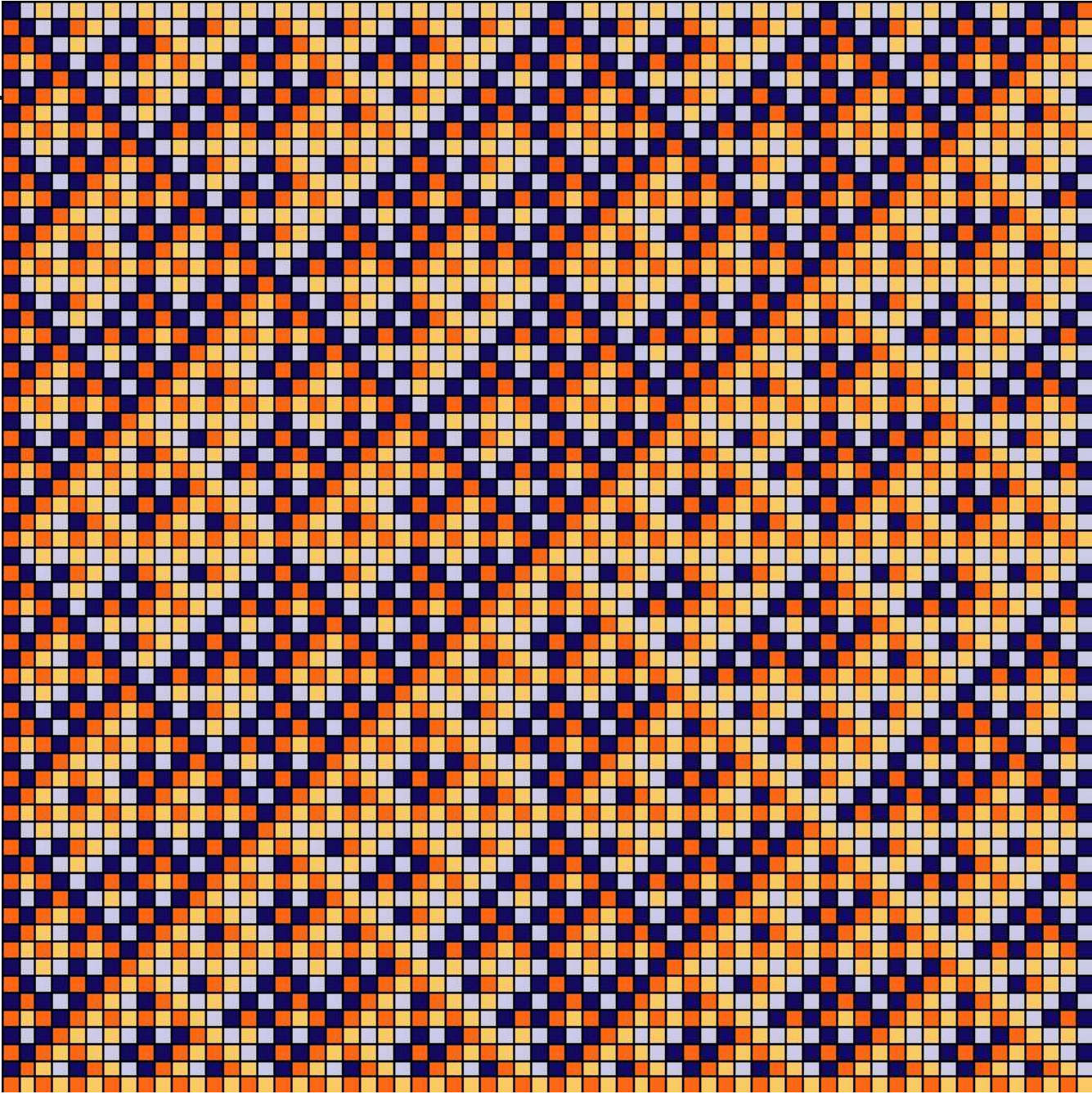
- boolean union, intersection, complement
- rotation
- *not* closed under concatenation, iteration
'regular'

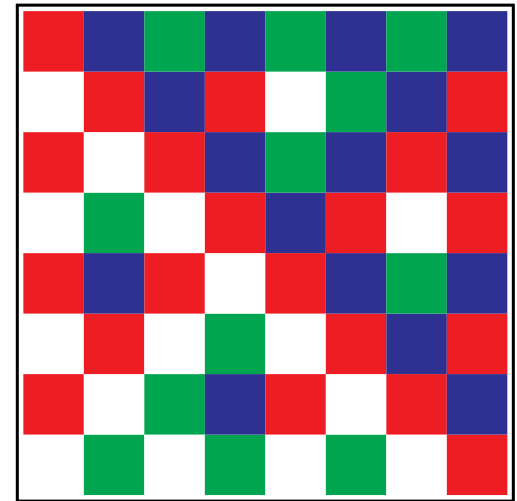
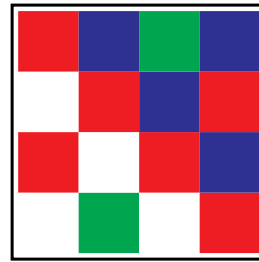
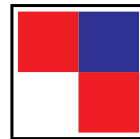
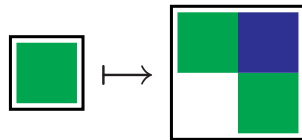
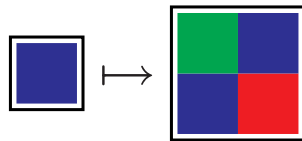
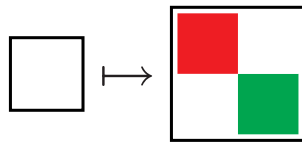
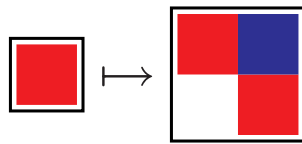
non-deterministic closed under

- union, intersection,
- rotation
- *not* closed under concatenation, iteration,
complement

4DFA \subset 4NFA \subset WANG

■ Iterated substitutions





square chair tiling

grammar-like

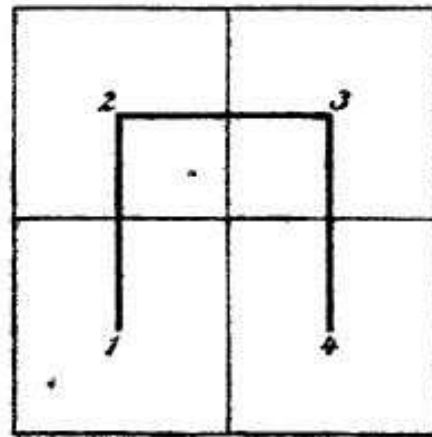


Fig. 1.

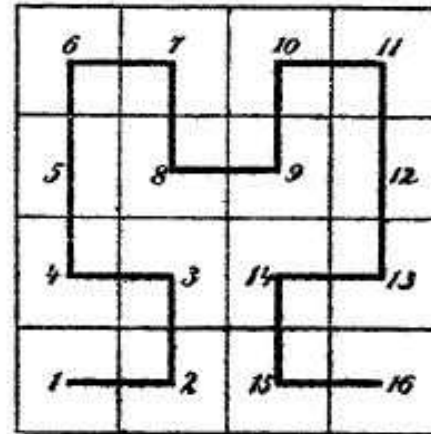


Fig. 2.

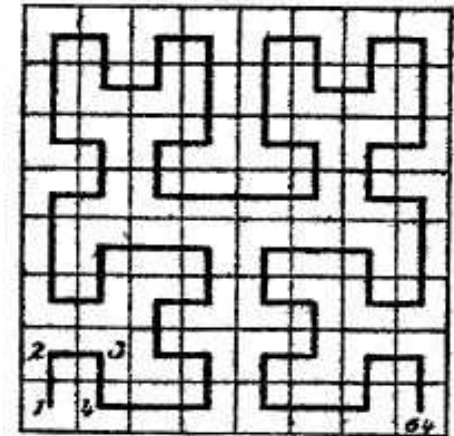
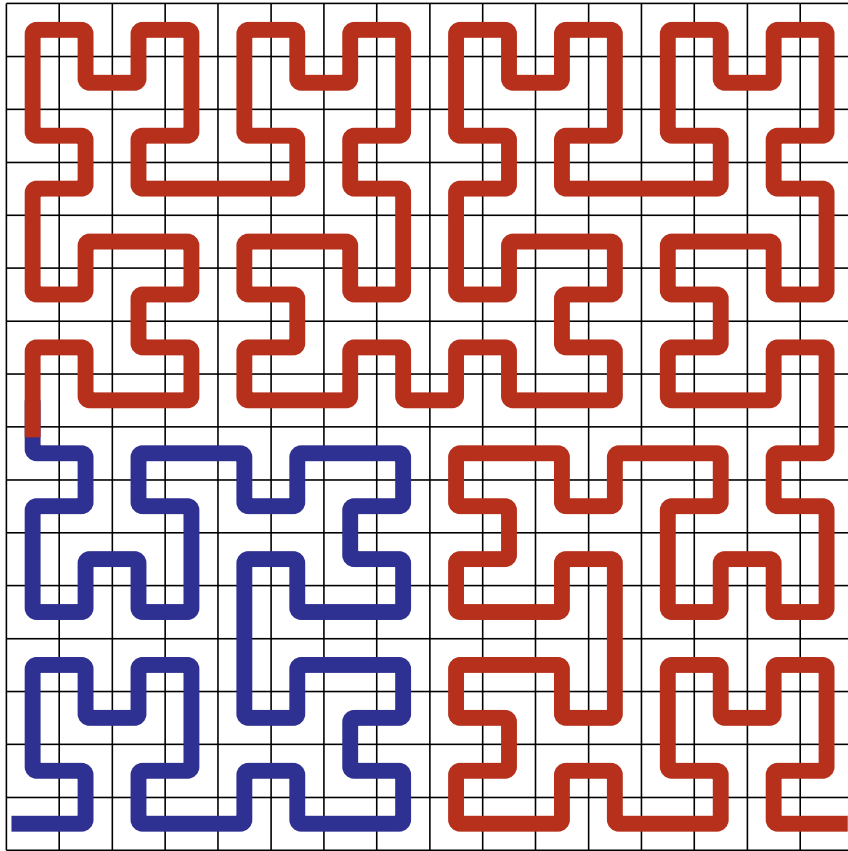


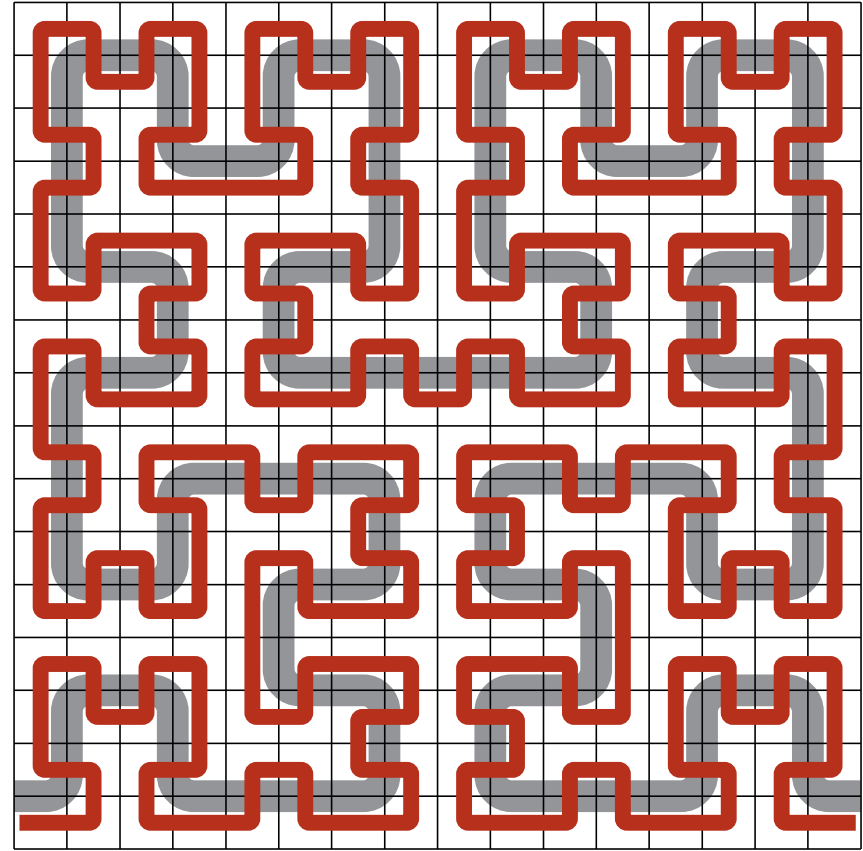
Fig. 3.

Ueber die stetige Abbildung einer Linie auf ein Flächenstück (1891)

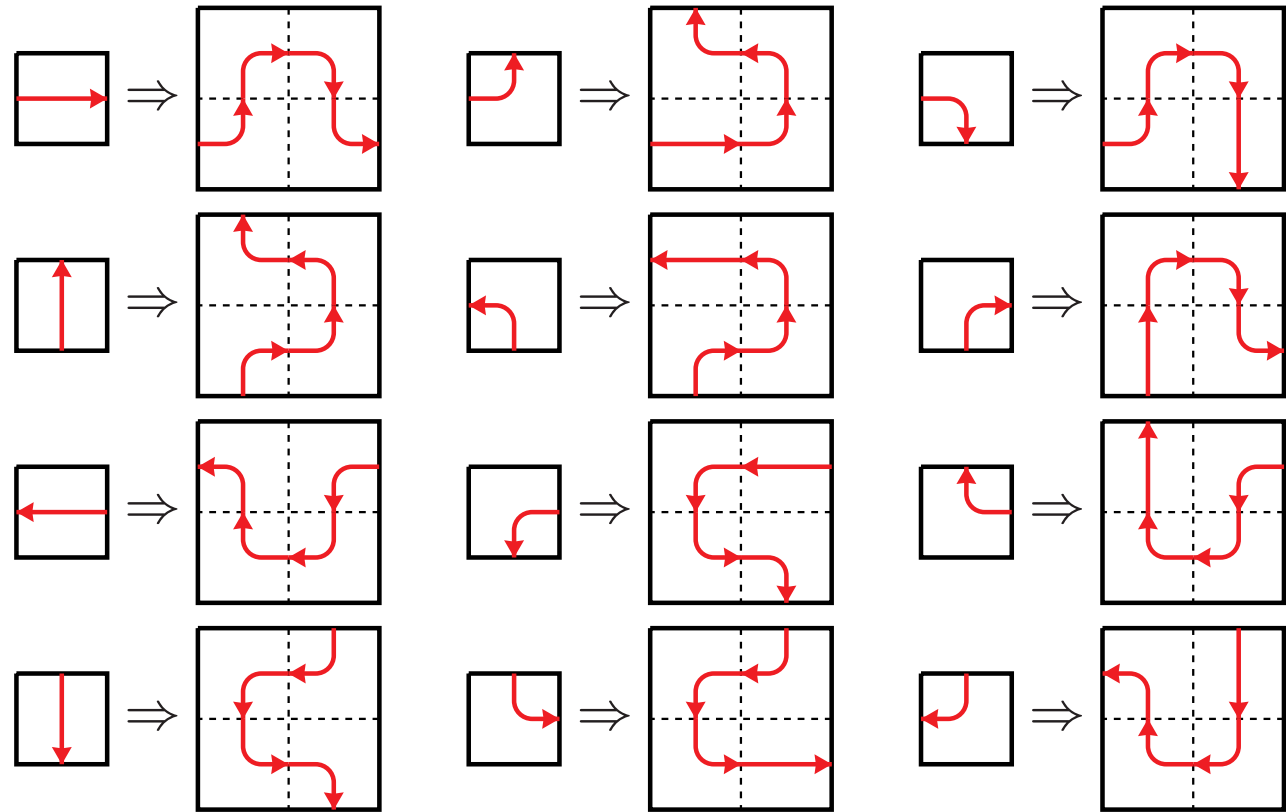
David Hilbert in Königsberg i. Pr.



four copies
top



refinement
bottom

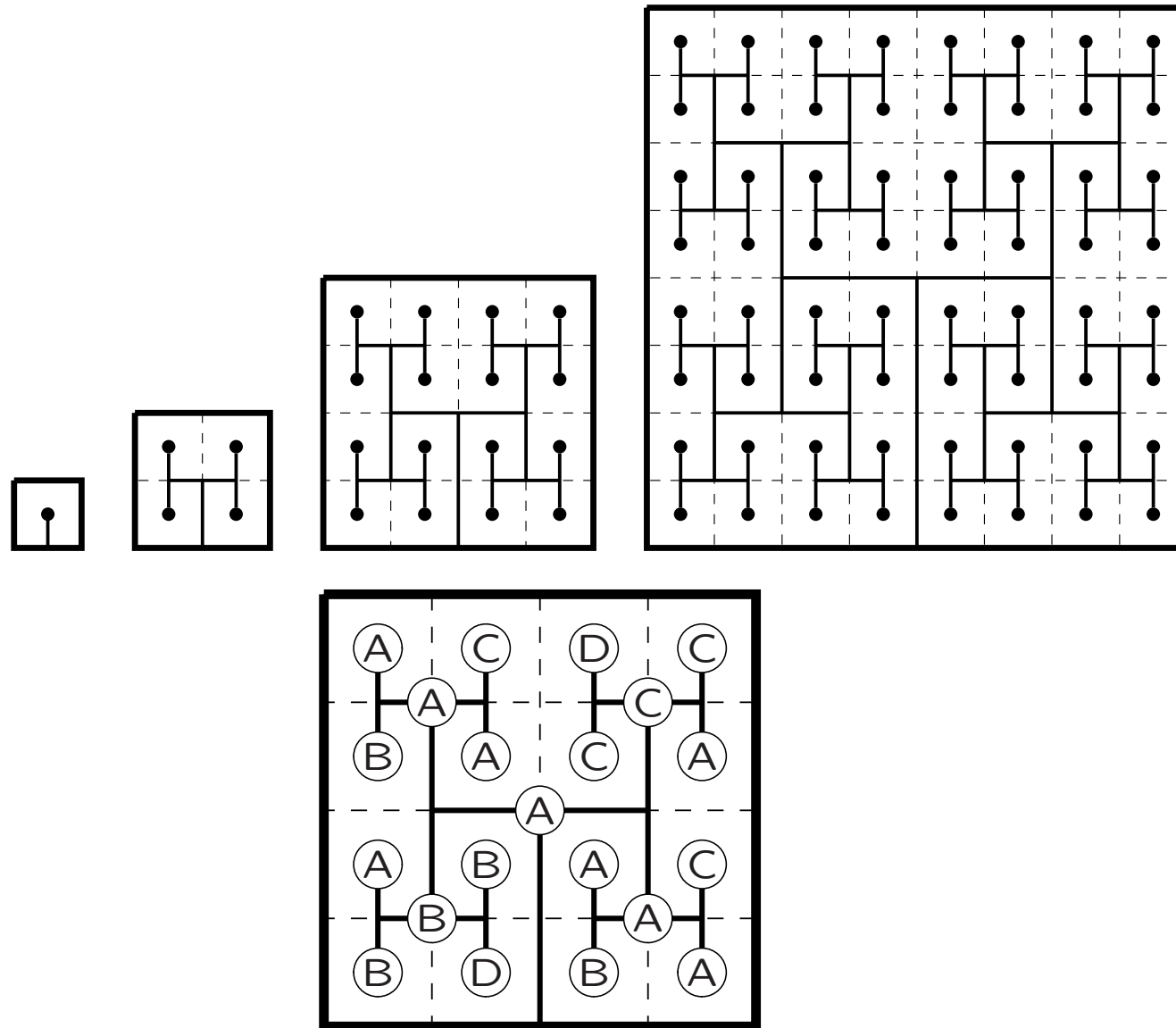


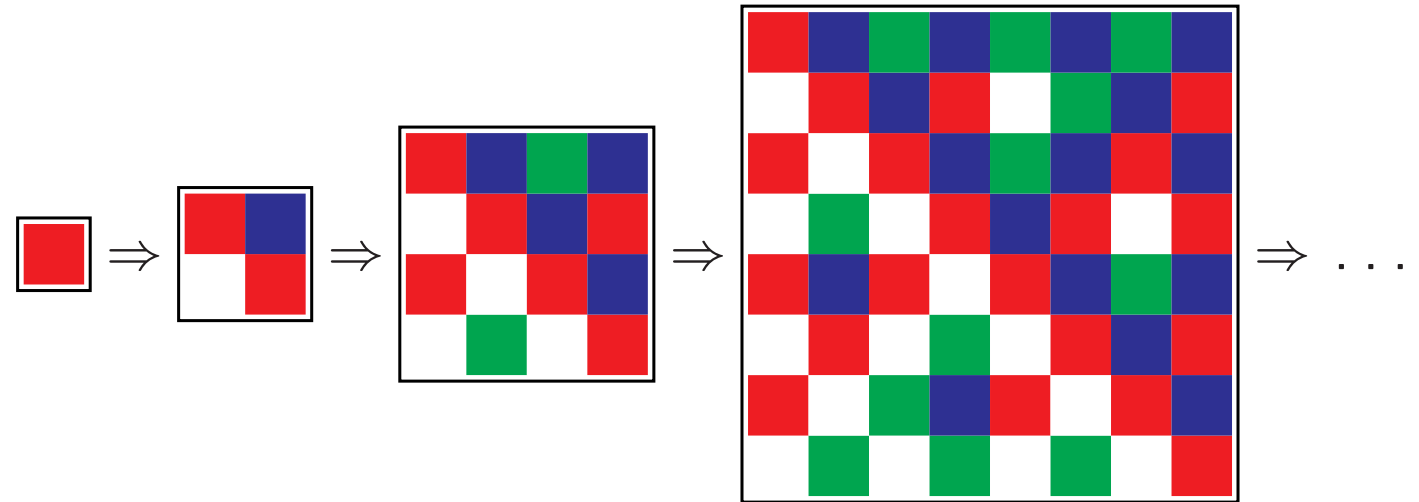
twelve tiles

Proposition. Every picture language defined by a substitution can be Wang tiled

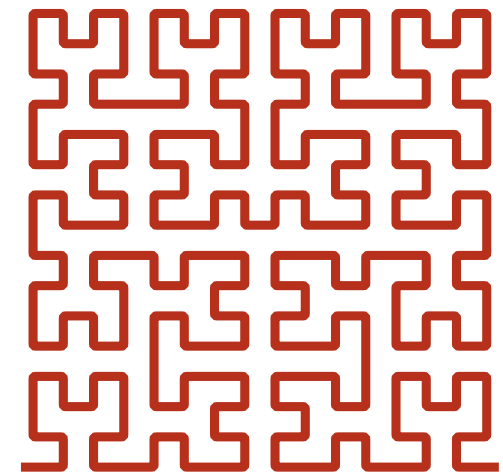
Lindgren, Moore, Nordahl: Complexity of two-dimensional patterns (1998)

Mozes: Tilings, Substitution Systems and Dynamical Systems Generated by Them (1989).

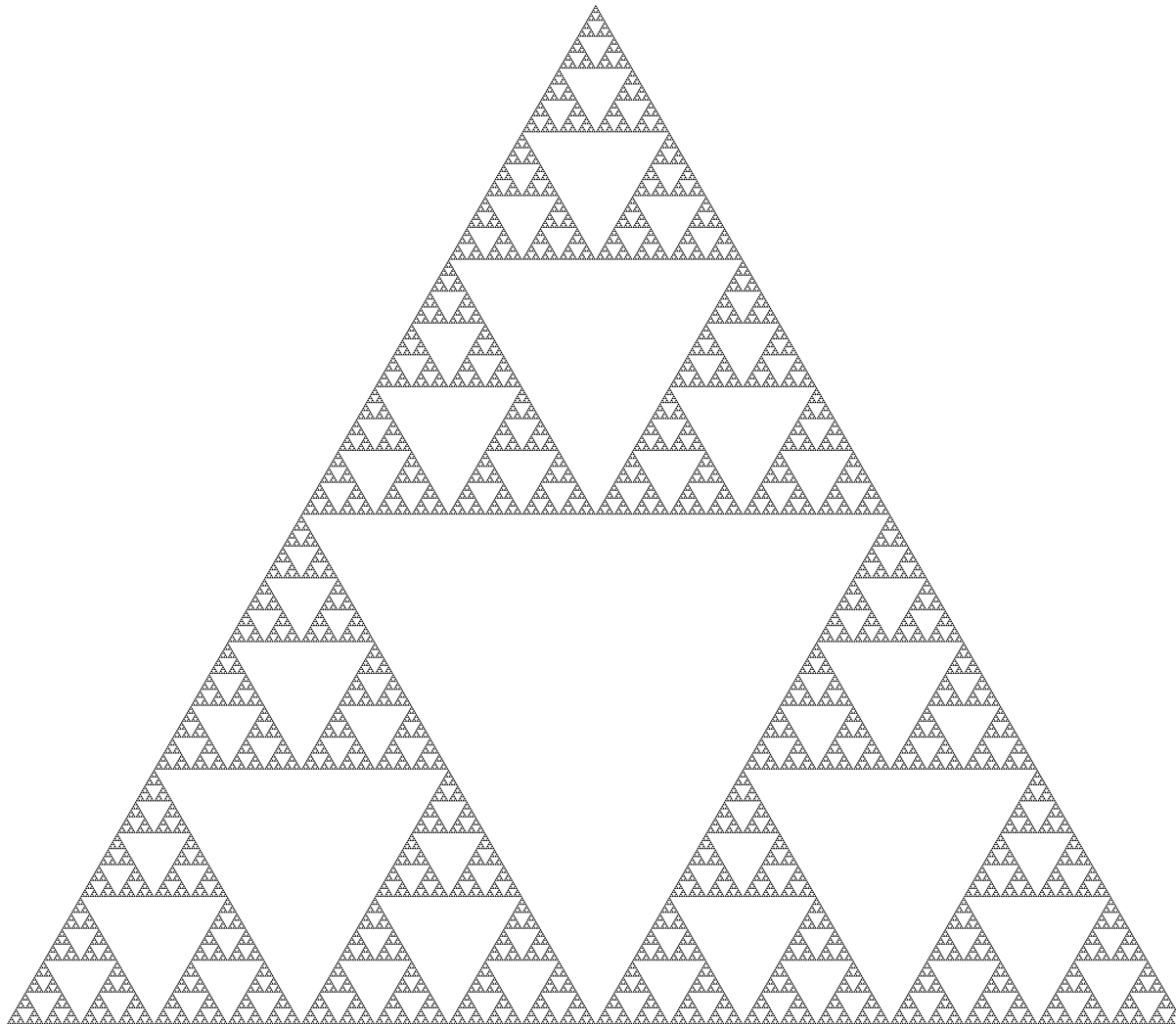




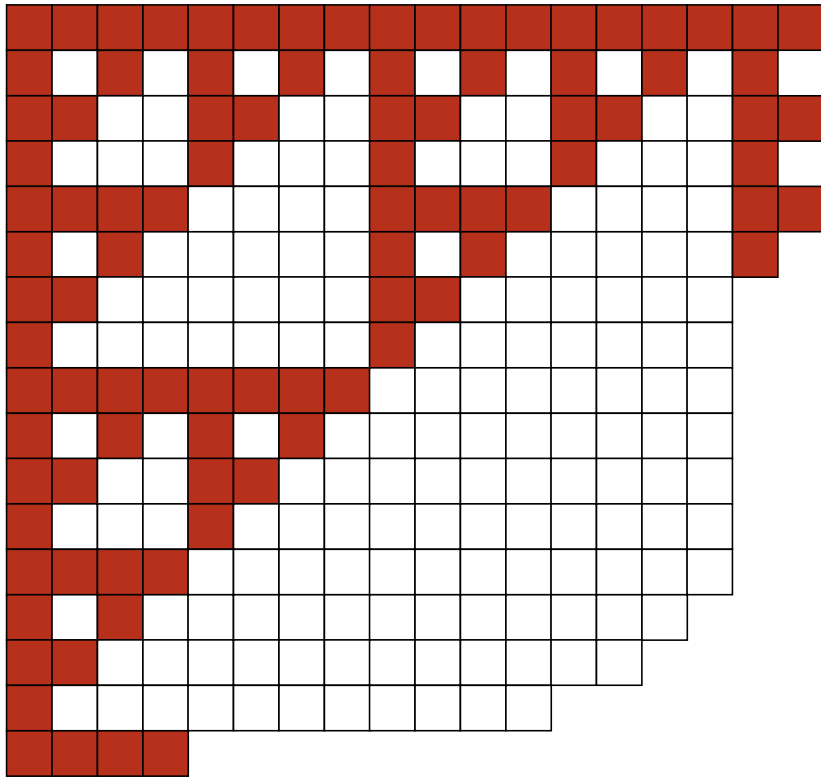
iterated substitutions form patterns that can be tiled



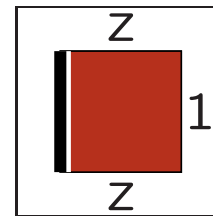
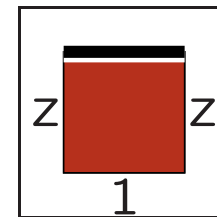
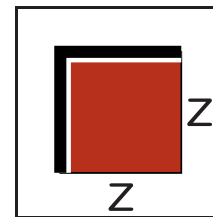
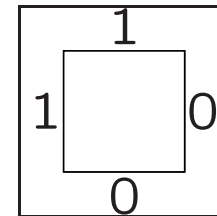
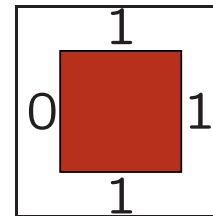
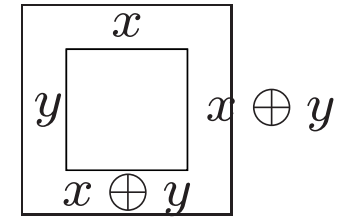
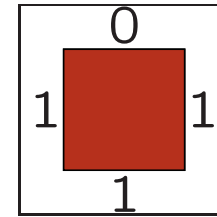
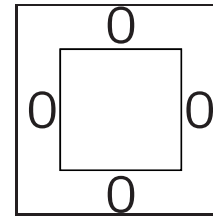
■ Self-Assembly

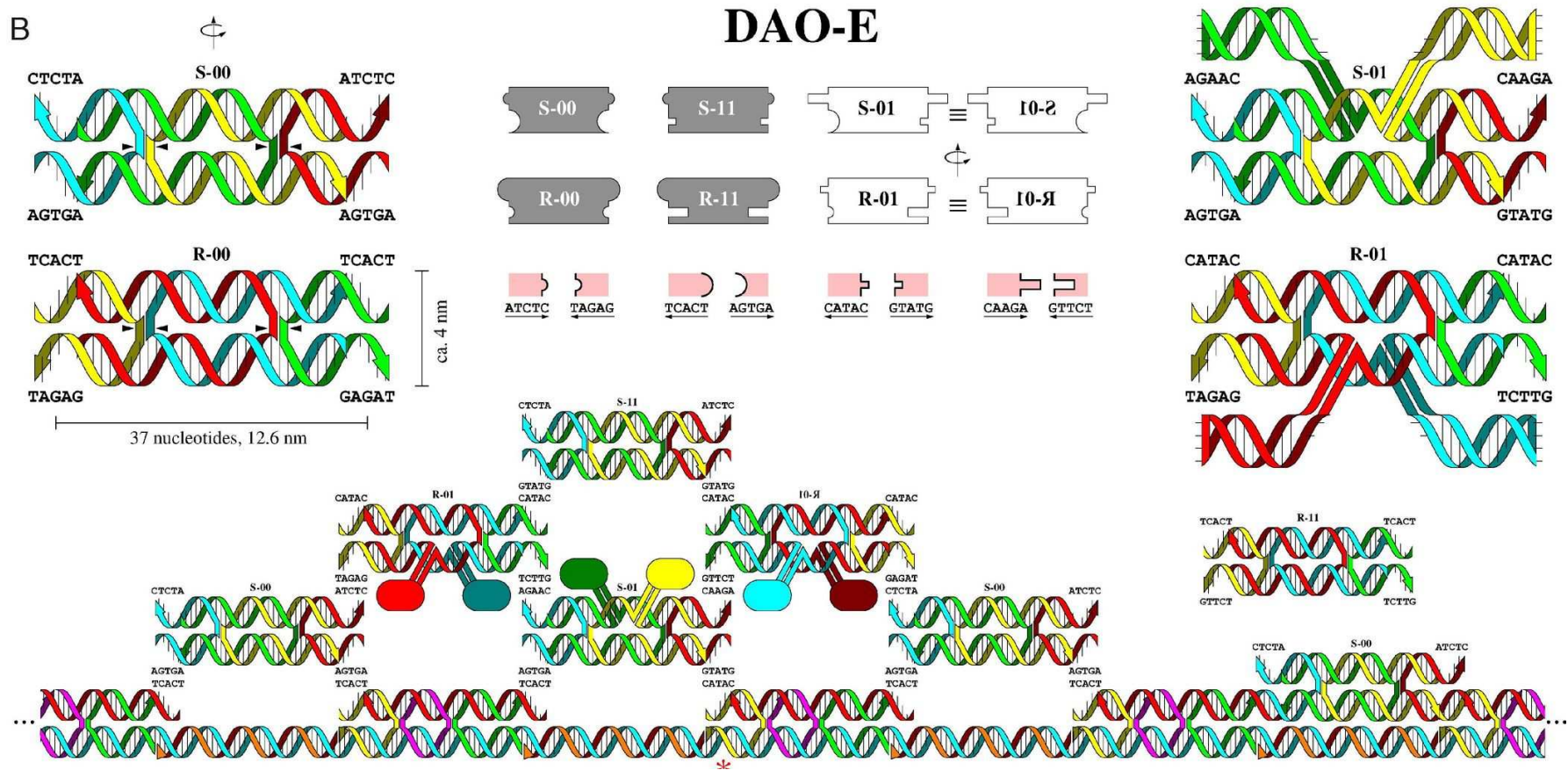


Qef's Website
wikipedia



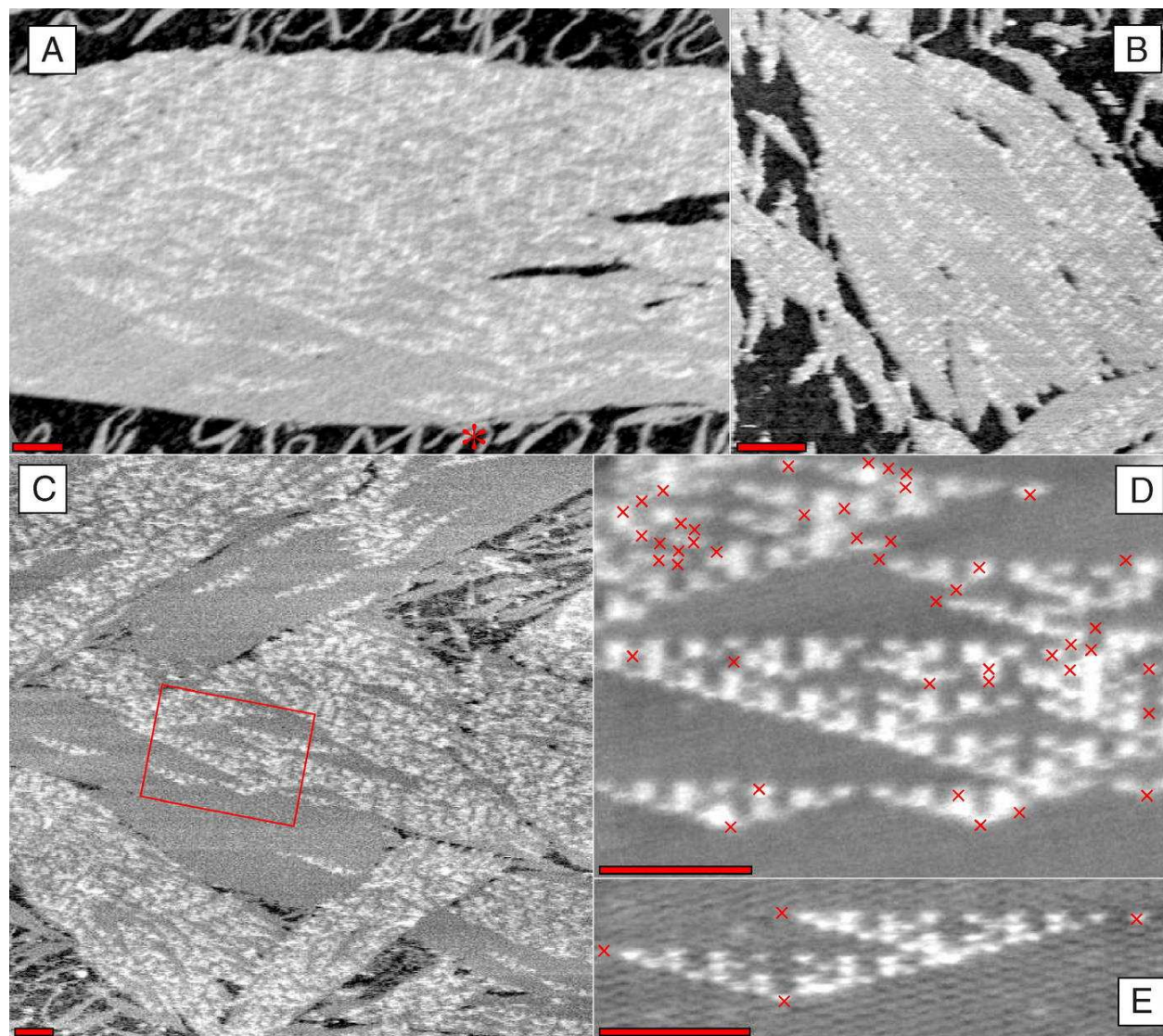
x	y	$x \oplus y$
0	0	0
0	1	1
1	0	1
1	1	0



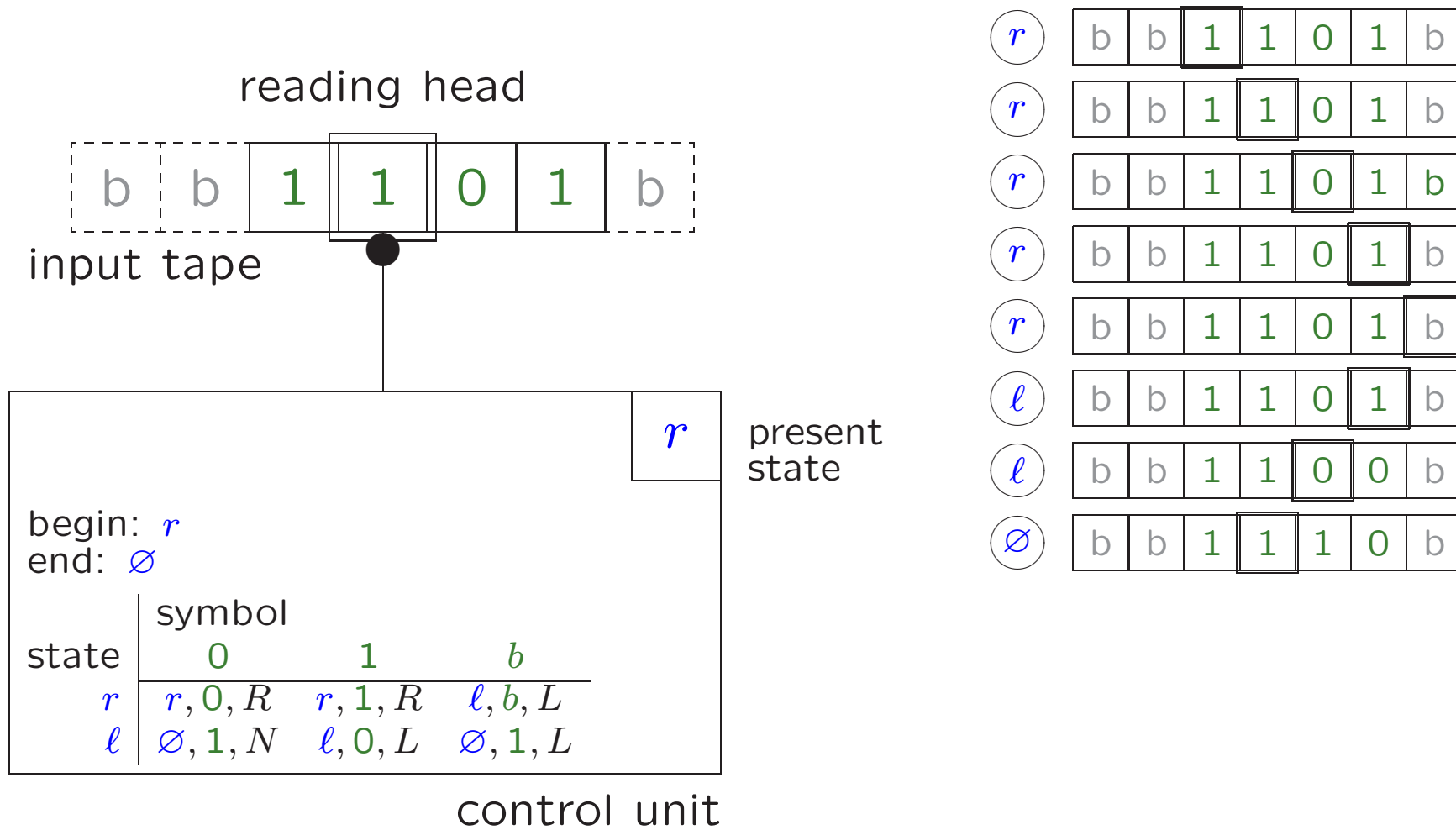


Algorithmic Self-Assembly of DNA Sierpinski Triangles (2004)

Rothemund, Papadakis, Winfree; PLoS Biology



■ no algorithm

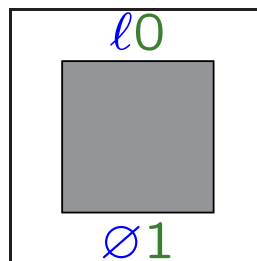
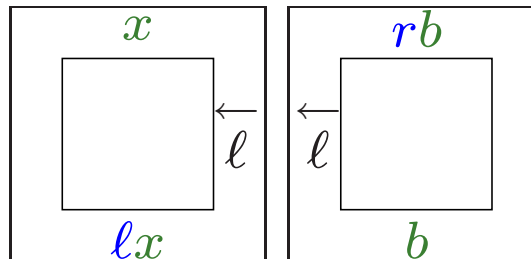
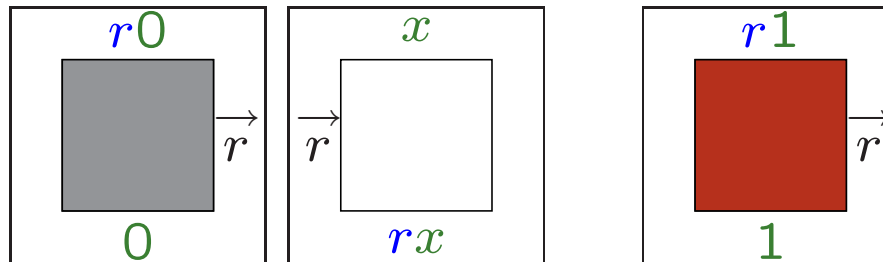
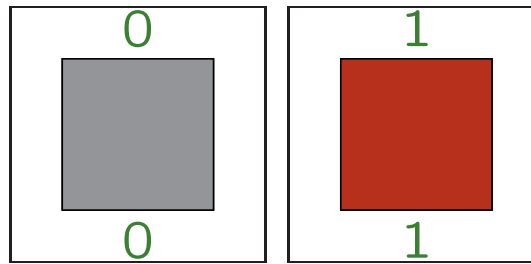


Turing, A. M. *On Computable Numbers, with an Application to the Entscheidungsproblem*. Proc. London Math. Soc. Ser. 2 42, 230-265, 1937.

<i>r</i> 0	1	0	1	1	<i>B</i>
0	<i>r</i> 1	0	1	1	<i>B</i>
0	1	<i>r</i> 0	1	1	<i>B</i>
0	1	0	<i>r</i> 1	1	<i>B</i>
0	1	0	1	<i>r</i> 1	<i>B</i>
0	1	0	1	1	<i>rB</i>
0	1	0	1	<i>l</i> 1	<i>B</i>
0	1	0	<i>l</i> 1	0	<i>B</i>
0	1	<i>l</i> 0	0	0	<i>B</i>
0	1	1	0	0	<i>B</i>



	0	1	b
r	$r, 0, R$	$r, 1, R$	l, b, L
l	$\emptyset, 1, N$	$l, 0, L$	$\emptyset, 1, L$



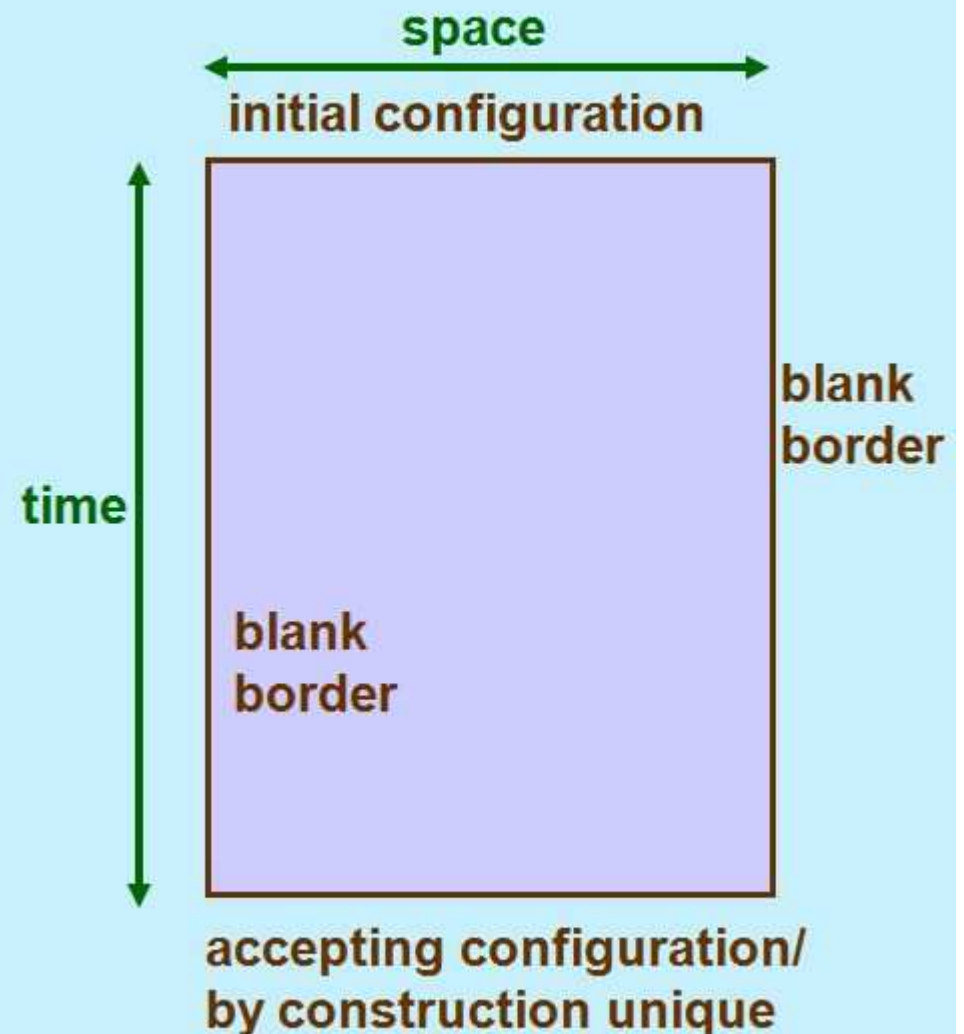
Tiling reductions

Program : Tile Types

Input: Boundary
condition

Space: Width region

Time: Height region

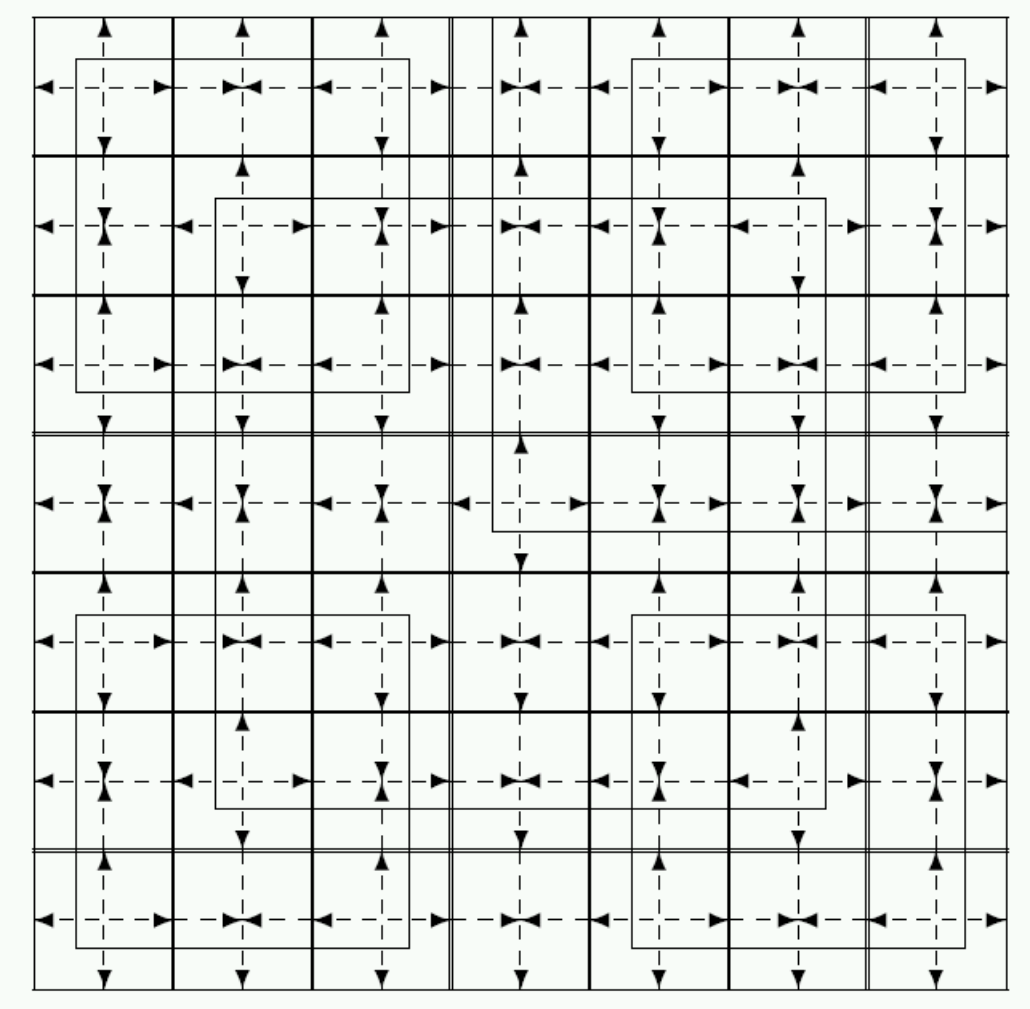


Square Tiling: A given square with boundary condition:
Complete for NP.

Corridor Tiling: A rectangle with boundary conditions on entrance and exit (length is undetermined):
Complete for PSPACE .

Origin Constrained: The entire plane with a given Tile at the Origin. **Complete for co-RE**
hence **Undecidable**

General: The entire plane without constraints. Still **Complete for co-RE** (Wang/Berger's Theorem). **Hard to Prove!**



Robinson

Bedankt!

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